

High-speed high-security
signatures

D. J. Bernstein

University of Illinois at Chicago

Niels Duif

Technische Universiteit Eindhoven

Tanja Lange

Technische Universiteit Eindhoven

Peter Schwabe

National Taiwan University

Bo-Yin Yang

Academia Sinica

Fast signature verification.

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280880 cycles for verification.

(Official **eBATS** measurement.)

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Very fast signing.

88328 cycles for signing.
(Official eBATS measurement.)
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64 bytes for signature.

No hidden slowdowns.

Small keys.

32 bytes for public key.

No hidden slowdowns.

No secret array indices.

No information flow from
secret data to addresses.

⇒ No cache-timing attacks.

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No secret branch conditions.

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Software uses the
NaCl/SUPERCOP API:
`crypto_sign_keypair`,
`crypto_sign`,
`crypto_sign_open`.

Available now in SUPERCOP:
bench.cr.yp.to/supercop.html

Will also be in NaCl:
nacl.cr.yp.to

Public domain—
use it any way you want!

Paper is online:
ed25519.cr.yp.to