

ChaCha,  
a variant of Salsa20

D. J. Bernstein

University of Illinois at Chicago

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Note: ChaCha is

not an eSTREAM candidate!

“Post-eSTREAM cryptography.”

## The Salsa20 cipher family

Salsa20/20 is faster than AES.

I recommend Salsa20/20

as an AES replacement

for typical applications.

What if the users value speed

more highly than confidence?

Reduce rounds: Salsa20/12 or /8.

Attacks: Crowley, SASC 2006;

Fischer et al., Indocrypt 2006;

Tsunoo et al., SASC 2007;

Aumasson et al., FSE 2008.

Against Salsa20/8, slightly faster

than 256-bit exhaustive search.

## The ChaCha cipher family

ChaCha8, ChaCha12, ChaCha20  
are like Salsa20/8 etc.:

same key size, same block size,  
same arithmetic operations.

But there are some changes  
in the order of operations.

Aumasson et al. attack  
a preliminary “ChaCha”  
with one of these changes.

The change makes the attack  
fail one round sooner.

Perhaps the other changes  
are also helpful. Need analysis!

Changed order of operations  
often saves some CPU time.

Cycles/byte, 576-byte packet,  
for ChaCha8 and Salsa20/8:

C8	S8	
2.00	2.14	ppc32 PowerPC G4
2.09	2.07	amd64 Core 2 Duo
3.47	3.66	amd64 Athlon 64 X2
4.06	5.57	amd64 Pentium D f64
5.05	5.54	x86 Pentium M 695
6.11	5.98	x86 Pentium 4 f12
6.27	7.06	x86 Pentium 4 f41
6.71	6.76	sparc UltraSPARC III
7.02	7.20	ppc64 Cell PPE

## Salsa20 quarter-round

4 modifications per quarter-round:

$$x_4 \hat{=} (x_0 + x_{12}) \lll 7$$

$$x_8 \hat{=} (x_4 + x_0) \lll 9$$

$$x_{12} \hat{=} (x_8 + x_4) \lll 13$$

$$x_0 \hat{=} (x_{12} + x_8) \lll 18$$

16 modifications per round.

128 modifications in Salsa20/8.

320 modifications in Salsa20/20.

Key words kept separate? No!

(“I also see each use of a  $k$  word as a missed opportunity to spread changes through the  $n$  words.”)

Constants kept separate? No!

## ChaCha quarter-round

Salsa20 temporarily allocates separate word for storing sum.

Missed diffusion opportunity!

ChaCha takes this opportunity:

$x_0 += x_{12}$ ;  $x_4 \hat{=} x_0$ ;  $x_4 \lll = 16$

$x_8 += x_4$ ;  $x_{12} \hat{=} x_8$ ;  $x_{12} \lll = 12$

$x_0 += x_{12}$ ;  $x_4 \hat{=} x_0$ ;  $x_4 \lll = 8$

$x_8 += x_4$ ;  $x_{12} \hat{=} x_8$ ;  $x_{12} \lll = 7$

In absence of carries,

average 1-bit differential

affects 12.5 output bits.

Salsa20: only 8 output bits.

## Salsa20 round switch

Salsa20 sweeps down columns:

x4, x8, x12, x0;

x9, x13, x1, x5;

x14, x2, x6, x10;

x3, x7, x11, x15.

Then across rows:

x1, x2, x3, x0;

x6, x7, x4, x5;

x11, x8, x9, x10;

x12, x13, x14, x15.

Modifications in these two rounds:

below-diagonal, . . . , . . . , diagonal,

. . . , . . . , below-diagonal, diagonal.

## ChaCha round switch

ChaCha sweeps *up* columns:

x0, x12, x8, x4;

x5, x1, x13, x9;

x10, x6, x2, x14;

x15, x11, x7, x3.

Then across rows:

x0, x1, x2, x3;

x5, x6, x7, x4;

x10, x11, x8, x9;

x15, x12, x13, x14.

Diagonal, ..., ..., below-diagonal,  
diagonal, ..., ..., below-diagonal.

Intuition: better diffusion.



## Salsa20 input positions

Only 128 bits of input  
are controlled by attacker:

$x_6, x_7, x_8, x_9$ .

$2^{256}$  possible input pairs.

$\approx 2^{264}$  possible state pairs.

Wishful-thinking differentials  
impose  $\approx 512$  bit conditions  
and will never be entered.

See “Salsa20 security,” 2005.04.

Successful reduced-round attacks  
have traced 1-bit differential  
starting from “best” bit in  
 $x_6, x_7, x_8, x_9$ .

## ChaCha input positions

New choice of input positions:

$x_1, x_6, x_{11}, x_{12}$ .

All of these words are modified at second position in quarter-round.

Intuition: less variability

in “goodness” of these bits;

“best” bit for ChaCha isn’t

as good as “best” bit for Salsa20.

Also: Output array is permuted.

Better speed; identical security.

Convenient to describe states with the same permutation.

See paper for details.