

Differential addition chains

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Motivating problem:

Given elliptic curve E ,

integer n , and point P on E ,

compute nP on E

as quickly as possible.

Many variations of problem.

Some applications reuse one n

for many P 's.

Some applications don't.

Some applications use secret n ;

must not leak n through timing.

Some applications use public n .

Etc.

1987 Montgomery:

Focus on large-characteristic

curves $y^2 = x^3 + ax^2 + x$

with small $a \in \{6, 10, 14, \dots\}$.

Use pair (x, z) to represent point

$P = (x/z, \dots)$.

Computing $Q, R, Q - R \mapsto Q + R$

takes 6 mults.

Only 5 mults if $Q - R$ has small denominator.

Only 4 mults if $Q - R$ has small numerator and small denominator.

Only 4 mults if $Q = R$.

Given n , write $P \mapsto nP$

as composition of additions

$Q, R, Q - R \mapsto Q + R$.

e.g. $n = 10$: compute

$P, P, 0 \mapsto 2P$ with 4 mults;

$2P, P, P \mapsto 3P$ with 6 mults;

$3P, 2P, P \mapsto 5P$ with 6 mults;

$5P, 5P, 0 \mapsto 10P$ with 4 mults.

Overall 20 mults for $P \mapsto 10P$.

Only 18 mults

if P has small denominator.

Only 16 mults

if P has small numerator
and small denominator.

$0, P, 2P, 3P, 5P, 10P$ is a **differential addition chain** starting from $0, P$:
each subsequent term is $Q + R$ for some $Q, R, Q - R$ already in chain.

$0, 1, 2, 3, 5, 10$ is a differential addition chain starting from $0, 1$.

Question: Given n , how to find short differential addition chain starting from $0, 1$ and ending n ?

Variations: measure shortness by mults, CPU cycles, etc.

The binary method:

obtain $n, n + 1$ from

$\lfloor n/2 \rfloor, \lfloor n/2 \rfloor + 1$ using

one addition with difference 1,

one addition with difference 0.

e.g.

$13P, 13P, 0 \mapsto 26P$ with 4 mults;

$14P, 13P, P \mapsto 27P$ with 5 mults,

if P has small denominator.

Overall 9 mults

for each bit of n ,

if P has small denominator.

1992 Montgomery,

1996 Bleichenbacher,

2001 Tsuruoka: Can do better!

Experiments for average 128-bit n
find length ≈ 1.533 per bit,
instead of 2 per bit.

Lower bound ≈ 1.440 per bit.

Count mults instead of length:
 ≈ 8.885 per bit,
instead of 9 per bit.

Disadvantages: harder to find;
no uniform structure; harder to
avoid leaking n through timing.

Two-dimensional question:

Given m, n , how to find
short differential addition chain
starting from the vectors
 $(0, 0), (1, 0), (0, 1), (1, -1)$
and ending (m, n) ?

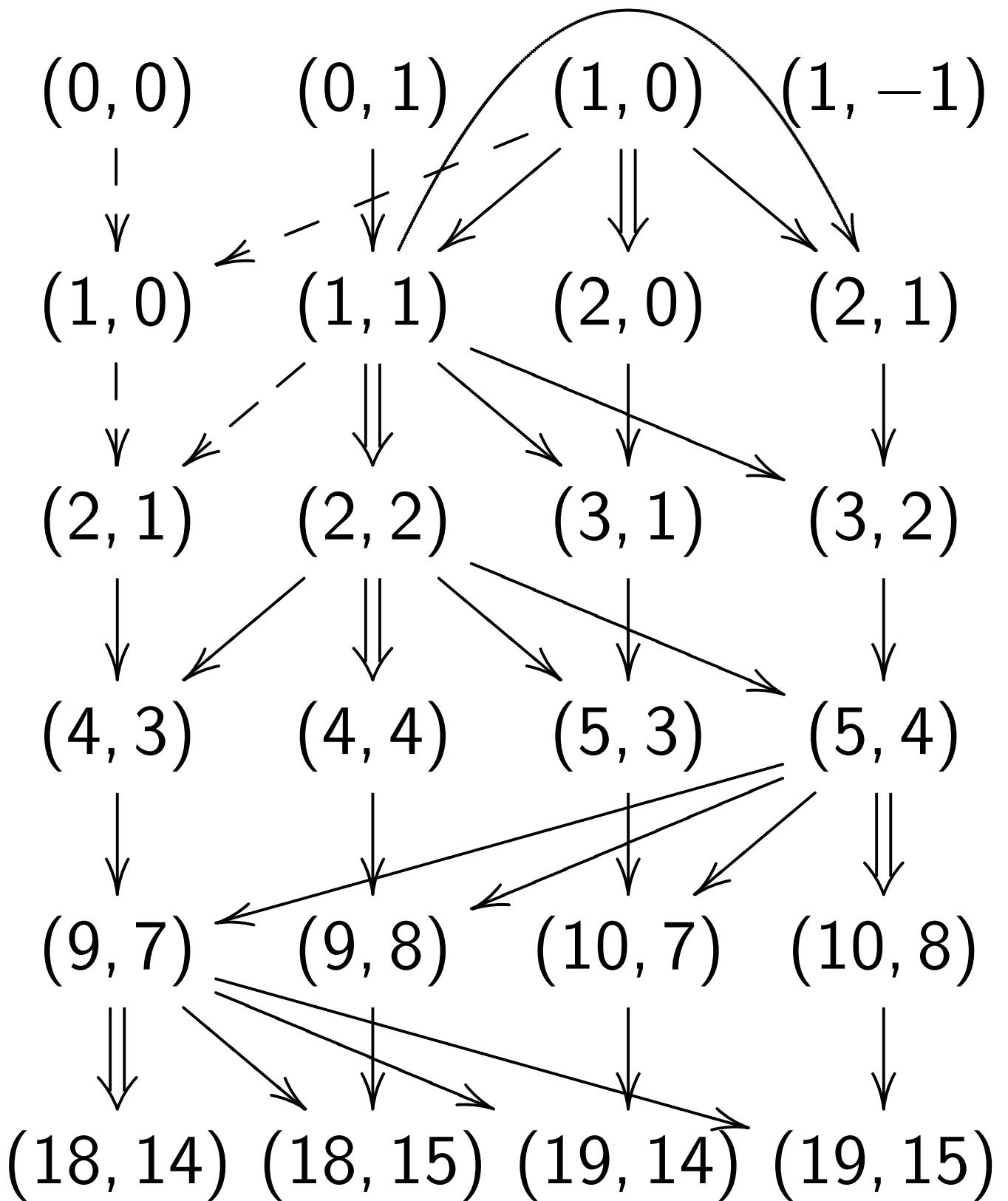
Motivating problem:

Given elliptic curve E ,
integers m, n ,
and points $P, Q, P - Q$,
compute $mP + nQ$ on E
as quickly as possible.

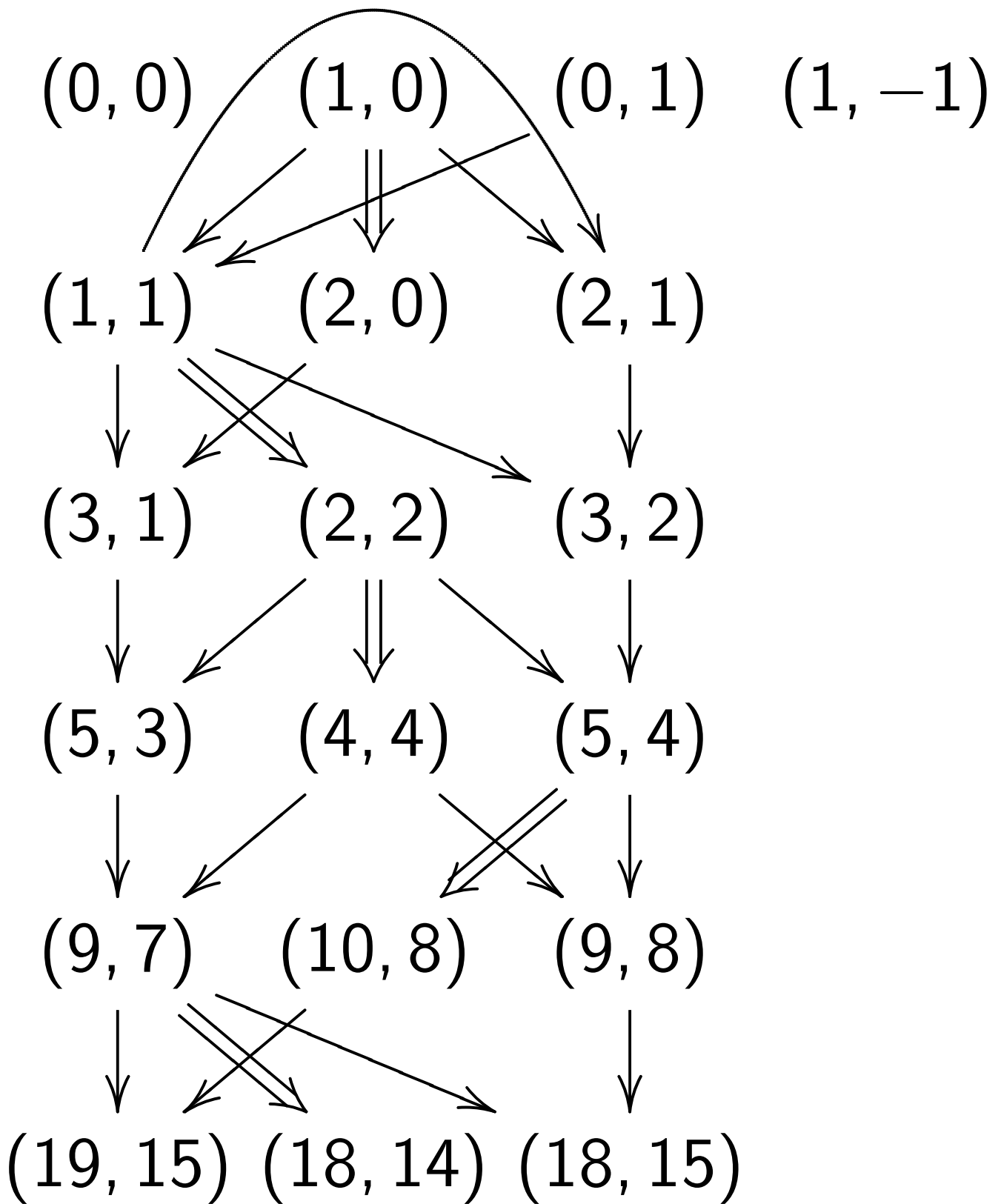
For average 128-bit exponents,
 small $P, Q, P - Q$ denominators:

dim	method	mults per bit	unif
2	easy binary	19.000	yes
2	Schoenmakers	17.250	no
2	Akishita	14.250	no
2	new binary	14.000	yes
2	Montgomery	10.261	no
2	new ext gcd	9.918	no
1	easy binary	9.000	yes
1	standard	8.885	no
	Fibonacci case	8.643	

Easy dim-2 binary chain:



New dim-2 binary chain:



Line in easy binary chain
has (a, b) , $(a, b + 1)$, $(a + 1, b)$,
 $(a + 1, b + 1)$. Obtain next line
by double-add-add-add.

New observation: can omit
(even, odd) or (odd, even),
chosen recursively so that
next line can be obtained
by double-add-add.

14 mults if $P, Q, P - Q$
have small denominators.

Intermediate results: 2000
Schoenmakers, 2001 Akishita.

How to do better than binary?

Don't worry about uniformity.

Critical idea for dim 1:

Build chain $0, 1, \dots, n$

by choosing $r \approx n(\sqrt{5} - 1)/2$

and building chain

$0, 1, \dots, r, n - r, n$.

Try many r 's, keep best.

Some further choices here:

could build $\{r, n - r, n\}$

from $\{r, n - 2r, n - r\}$ or

from $\{n - r, 2r - n, r\}$ or

from $\{r, n/2 - r, n/2\}$ or \dots

e.g. $n = 100$, $r = 39$:

Build chain

0, 1, 2, 3, 5, 7, 12, 17, 22, 39, 61, 100

by building $\{39, 61, 100\}$

from $\{22, 39, 61\}$ etc.

What about dim 2?

Obvious adaptation of idea:

Build chain $\dots, (m, n)$

by choosing (q, r)

and building chain

$\dots, (q, r), (m - q, n - r), (m, n)$.

e.g. Work backwards from
(314, 271) and (194, 167) to
(120, 104), then (74, 63), then
(46, 41), then (28, 22), then
(18, 19), then (10, 3), then (8, 16).

Hmmm, what's the endgame?

How to build short chain with
 $\{(8, 16), (10, 3), (18, 19)\}$?

Several plausible approaches,
but all of them scale badly.

Normally this construction
is abandoned.

New observation:

Simple endgames work well

if $rm - qn = \Delta$

with, e.g., $\Delta = \pm 2^a 3^b$.

Often find very good chains.

Easy to find (q, r)

given (m, n, Δ) :

standard ext-gcd computation.

What if (m, n) not coprime?

Great! Exploit factor.

Try many good choices

for (Δ, q, r) , keep best.

Example of new chain:

$(0, 0)$, $(1, 0)$, $(0, 1)$, $(1, -1)$,
 $(1, 1)$, $(1, 2)$, $(2, 3)$, $(3, 5)$,
 $(4, 7)$, $(5, 9)$, $(9, 16)$, $(14, 25)$,
 $(19, 34)$, $(33, 59)$, $(38, 68)$,
 $(66, 118)$, $(71, 127)$, $(61, 109)$,
 $(132, 236)$, $(203, 363)$, $(264, 472)$,
 $(325, 581)$, $(528, 944)$,
 $(731, 1307)$, $(1259, 2251)$,
 $(1787, 3195)$, $(2518, 4502)$,
 $(3249, 5809)$, $(5036, 9004)$,
 $(6823, 12199)$, $(10072, 18008)$,
 $(16895, 30207)$, $(26967, 48215)$.