

Reuters, 2004.11.09:

“Worm breaks speed record from discovery to life

“A new computer worm emerged on Tuesday which broke the speed record from the announcement of a security vulnerability in Microsoft’s Internet Explorer to a full-blown virus that spreads in the wild.

“The vulnerability was discovered and made public by two hackers with aliases ‘ned’ and ‘SkyLined’ on Friday, and only four days later a worm exploiting the weakness was developed and set loose, several virus-trackers reported.

“Microsoft said the worm is a variant of MyDoom and that it was investigating the threat the worm poses.

“Some anti-virus companies said the new worm was different from MyDoom because it spreads via weblinks and not e-mail attachments.

“ ‘People will receive an e-mail saying that their PayPal account has been credited or that they are invited to watch a webcam. When they click on the link, just by viewing a site it executes code and infects the computer,’ said technical consultant Graham Cluley at Sophos Anti-Virus.

“Microsoft was expected to issue its monthly batch of security patches later on Tuesday, but the company could not immediately say if a patch for the new worm would be part of it.

“However, the U.S. software giant said that consumers who had installed Service Pack 2 for Windows XP were at a reduced risk.

“The weakness in Internet Explorer is known as the IFRAME buffer overflow vulnerability.”

[AUS-CERT admits that future exploits may work under SP2.]

2004.11.15: Guest lecture
by Jon Solworth, Director,
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CS.

2004.11.17: Midterm 2,
focusing on setuid and related topics.

Assignment due 2004.11.22: read
textbook Chapter 4.

Attacker blocking permission bits

Each process has, in system data, **umask** (“file mode mask”).

Typical umask: 022.

Another typical umask: 077.

Any permission bit in umask is removed from new files.

e.g. `open("foo", O_CREAT, 0666)`

creates `foo` with permissions

0644 if umask is 022;

or 0600 if umask is 077.

Umask is preserved by `execve`.

Joe can run a `setuid` program with `umask` set to `0777`.

Files created by program then have permissions `000`: not readable, not writable, even to the file owner.

`root` can read and write anyway, but maybe program is `setuid` to something other than `root`.

Fix: Program sets its own `umask`.

System-specific setuid problems

OS designer adds system data and neglects to consider effect of data after setuid exec.

(Even worse: considers effect, and blames the setuid programs.)

e.g. FreeBSD allows two processes to share their signal actions.

FreeBSD bug fixed 2001.07.09:
shared signal actions weren't un-shared by `execve`.

Any user can take over any setuid program.

Another Sendmail example

Bug sort-of-fixed 1996.09.17:

```
a->q_uid = daemon_uid;
a->q_gid = daemon_gid;
pw = getpwnam(user);
if (pw != NULL) {
    a->q_uid = pw->pw_uid;
    a->q_gid = pw->pw_gid;
}
```

getpwnam() looks for
a uid and gid in /etc/passwd.

e.g. getpwnam("djb") returns
uid and gid 1001 if /etc/passwd
has djb:*:1001:1001:...

Context: Sendmail delivers messages to accounts such as djb.

`/home/djb/.forward` can specify a program to run for each message; Sendmail runs that program under djb's uid.

To figure out djb's uid, Sendmail calls `getpwnam("djb")`, which reads `/etc/passwd` and returns 1001.

Sendmail calls `setuid(1001)`.

Sendmail also delivers messages to aliases such as `postmaster`.

`/etc/aliases` can specify a program to run for each message; Sendmail runs that program under uid 1 (daemon).

Sendmail calls `getpwnam("postmaster")`, which doesn't find `postmaster` in `/etc/passwd`; returns 0.

Sendmail sees the 0 and calls `setuid(1)`.

Joe runs Sendmail, telling it to deliver a message to joe.

Sendmail looks in `/home/joe/.forward`, which says `"Run /home/joe/evil."`

Oops, system is very busy.

Sendmail saves message in queue, along with the following note:

`"Deliver message to joe by running /home/joe/evil."`

Joe starts Sendmail again,
telling it to run the queue:

```
joe% sendmail -q
```

System is no longer busy.

Sendmail tries to deliver message
by running `/home/joe/evil`.

But what uid should it use?

Sendmail calls `getpwnam("joe")`
to find the uid and gid.

By setting resource limits,
Joe can make `getpwnam()` fail.
Easiest: file-descriptor limits.

`getpwnam()` returns 0,
even though `joe` is in `/etc/passwd`.

Sendmail runs `/home/joe/evil`
as uid 1.

Joe can now read and destroy
subsequent mailing-list deliveries.

Sendmail “fix”:

Remove file-descriptor limits.

But Joe can still force
`getpwnam()` to fail.

System has limit on
total number of open files
across all processes.

If Joe opens many files,
`getpwnam()` can't open more.

Joe can attack any program,
not just setuid programs,
in this way.

Underlying source of problem:
getpwnam() returns 0
for “permanent” errors
(user not in /etc/passwd)
and for “temporary” errors
(unable to open /etc/passwd).

For temporary errors,
Sendmail needs to try again later;
but Sendmail can't tell
whether the error was temporary.

For comparison: If open()
fails because file doesn't exist,
it sets errno to ENOENT.
If it fails because of fd rlimit,
it sets errno to EMFILE.
getpwnam() should use ESRCH.