

Shapes of Elliptic Curves

Daniel J. Bernstein

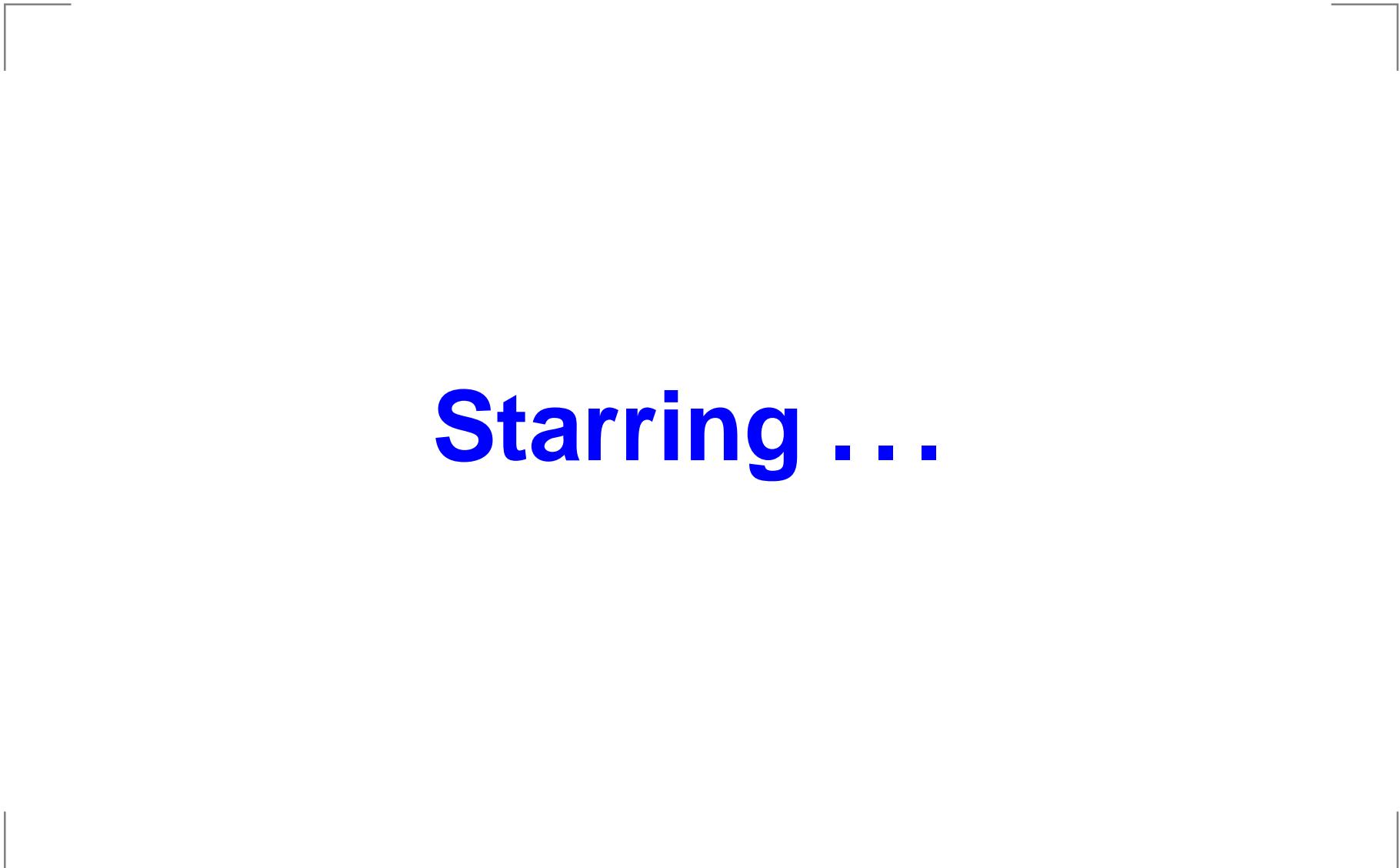
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19.05.2008



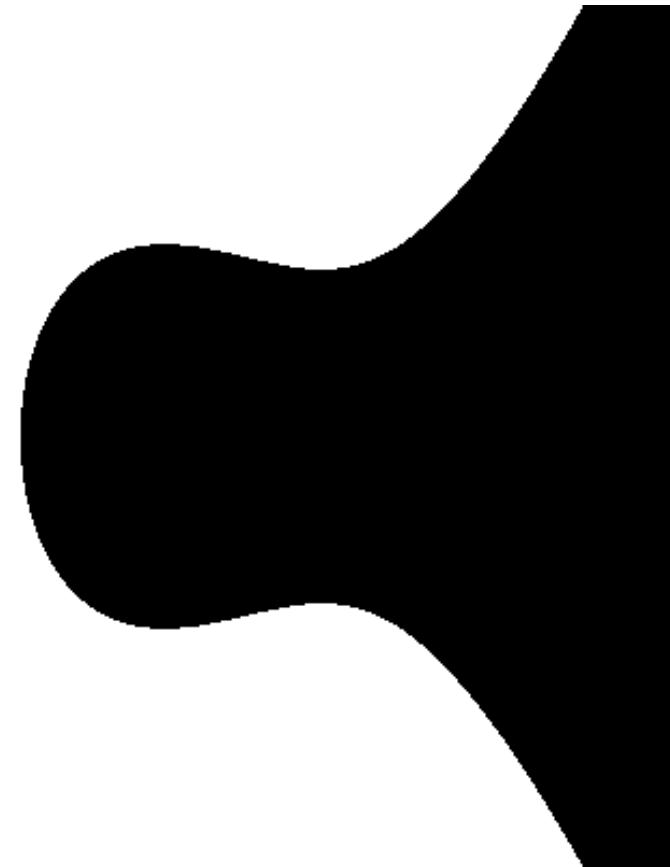
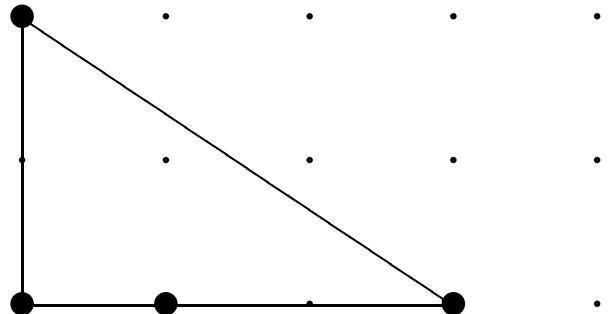
Starring . . .

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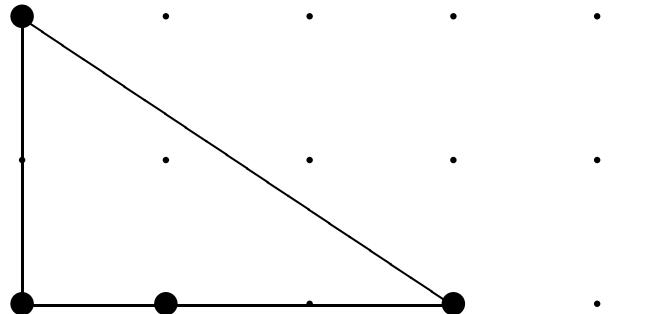
Weierstrass curve

$$y^2 = x^3 - 0.4x + 0.7$$



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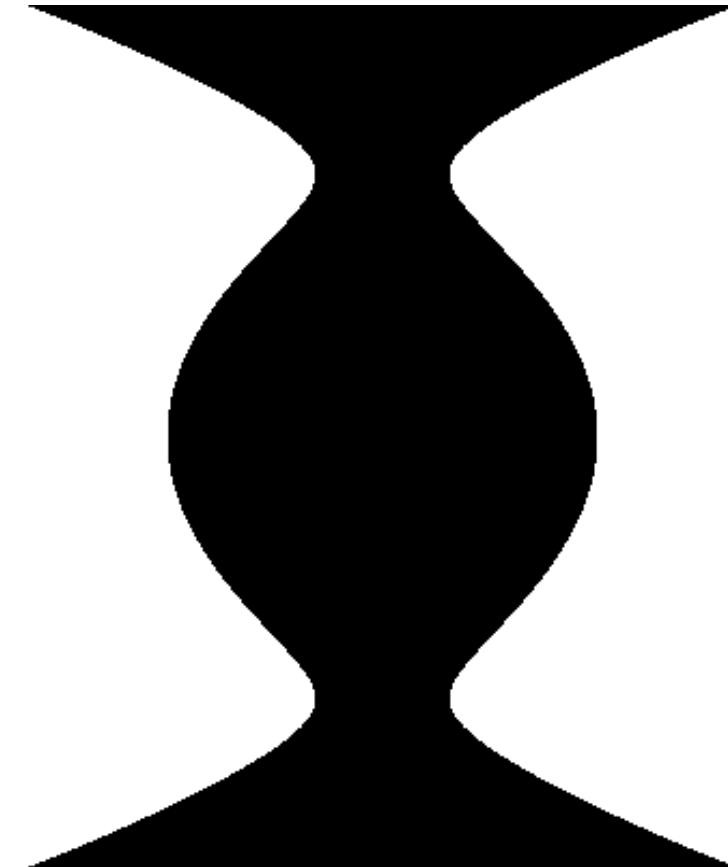
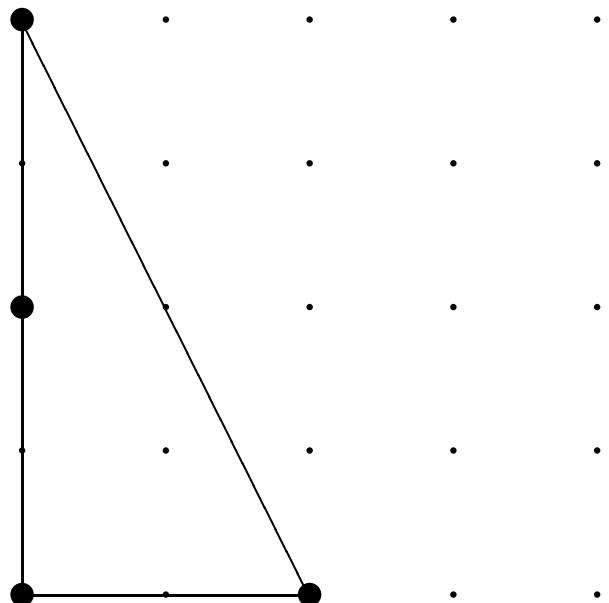


The Weierstrass-turtle: old, trusted and slow. Warning: (picture) incomplete!



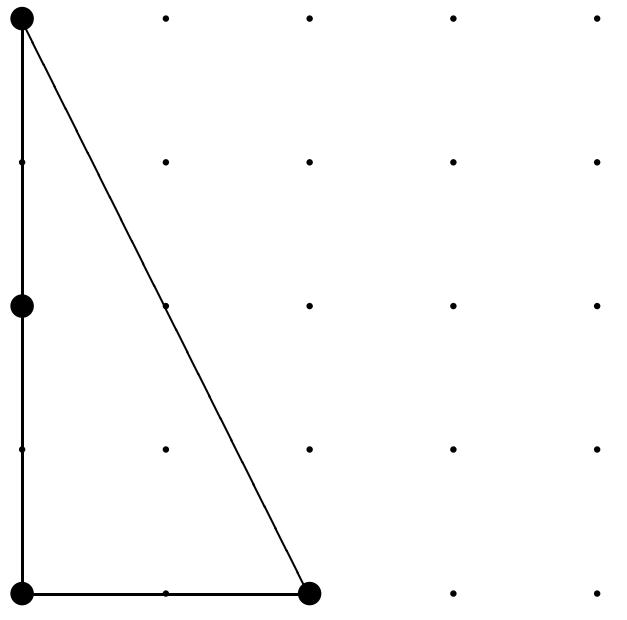
Jacobi quartic

$$x^2 = y^4 - 1.9y^2 + 1$$

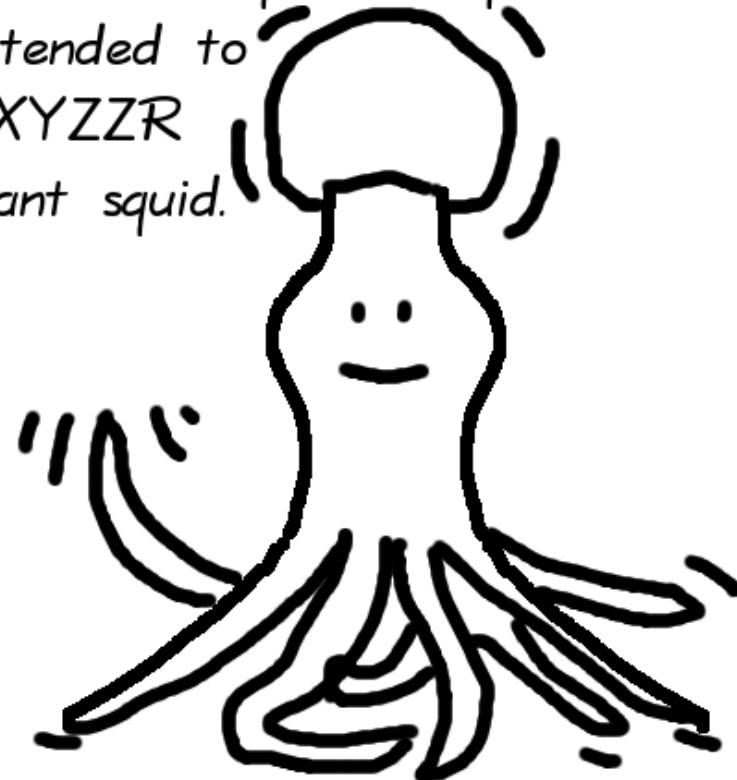


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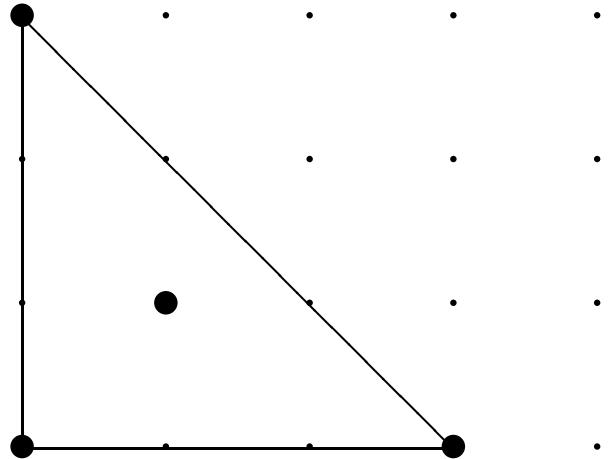


The Jacobi-quartic squid: can be
extended to
 $XXYZZR$
giant squid.



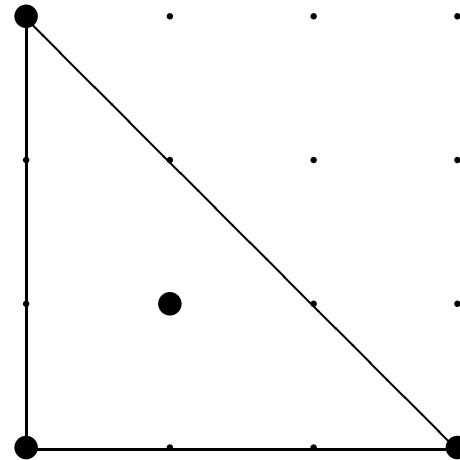
Hessian curve

$$x^3 - y^3 + 1 = 0.3xy$$



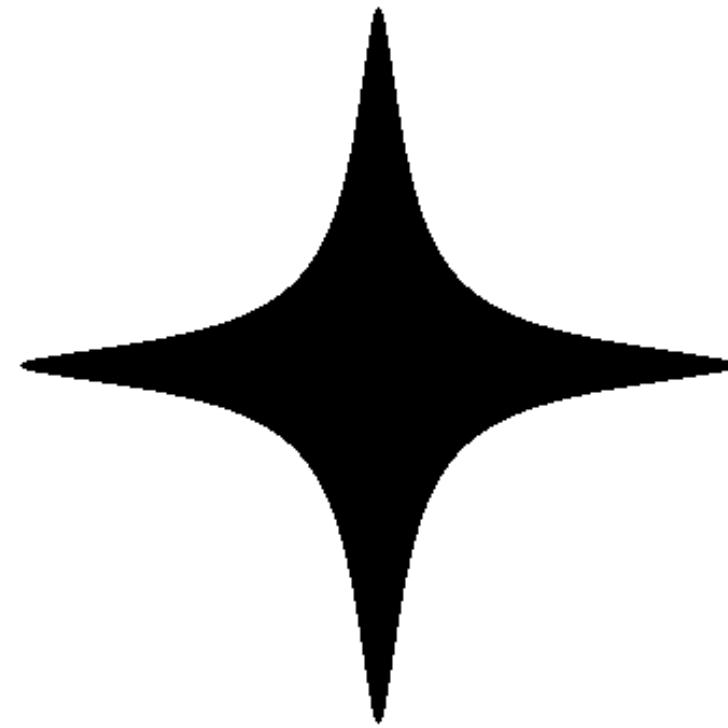
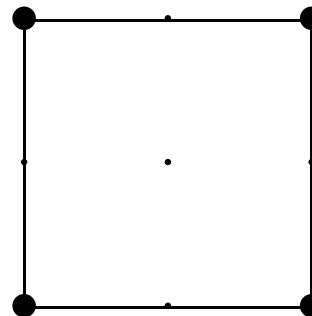
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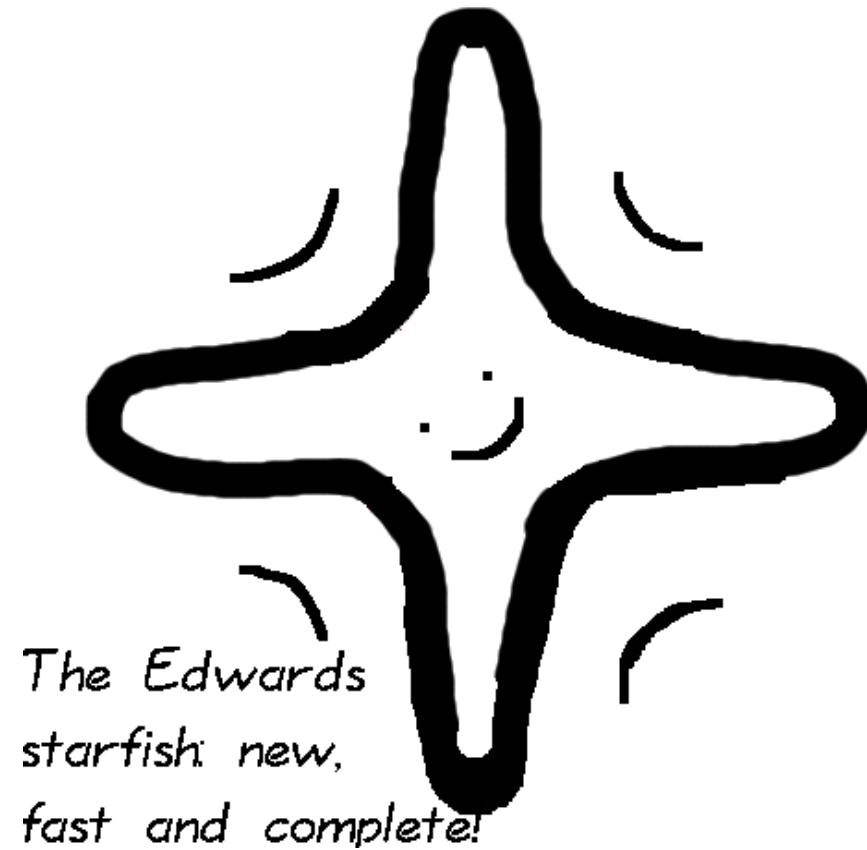
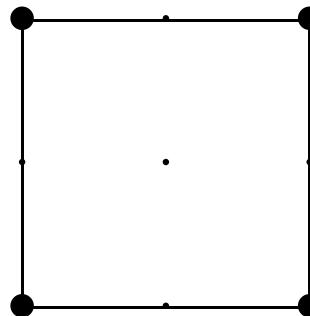
Edwards curve

$$x^2 + y^2 = 1 - 300x^2y^2$$



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The race – zoom on Weierstrass and Edwards

Weierstrass vs. Edwards I

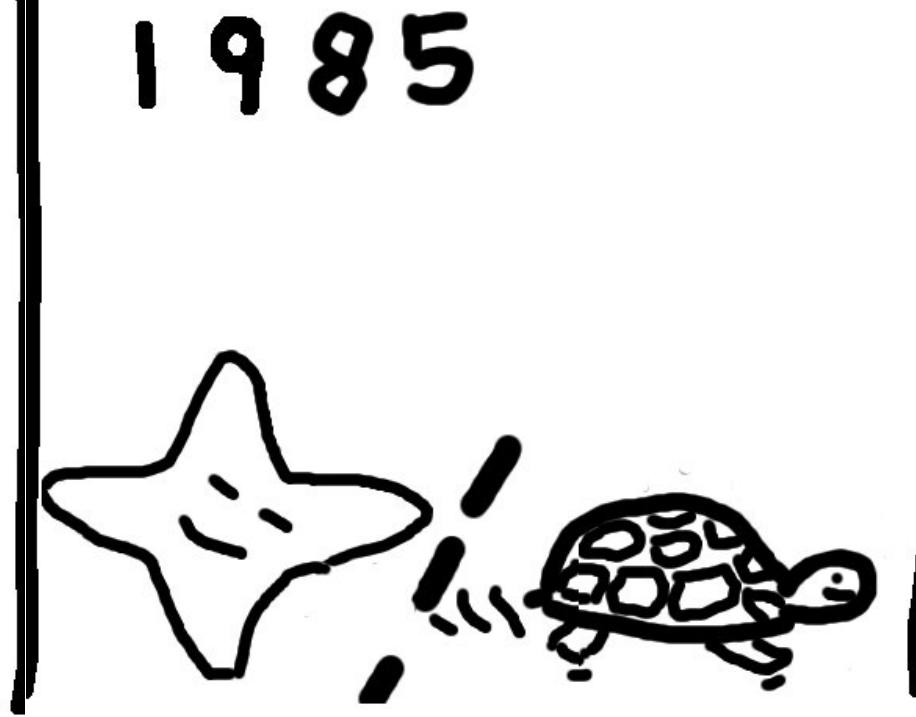


Start!

Weierstrass vs. Edwards II



Start!



Weierstrass sets off, Edwards
left behind sleeping

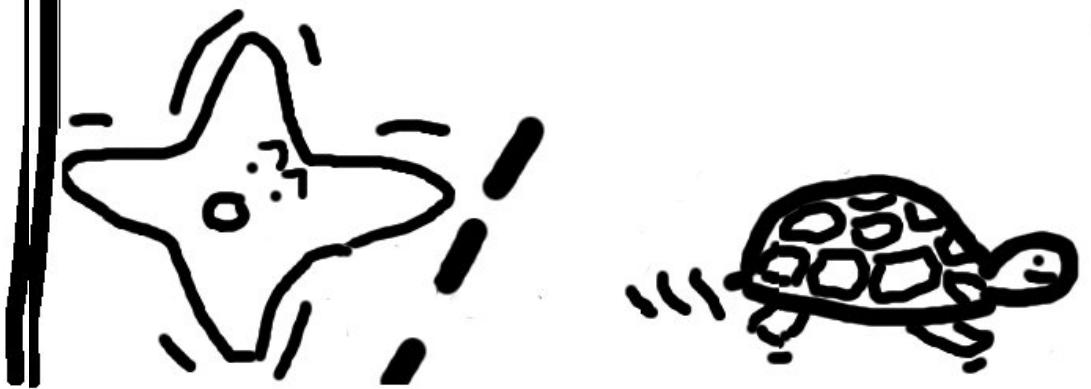
Weierstrass vs. Edwards III

1985



Weierstrass sets off, Edwards
left behind sleeping

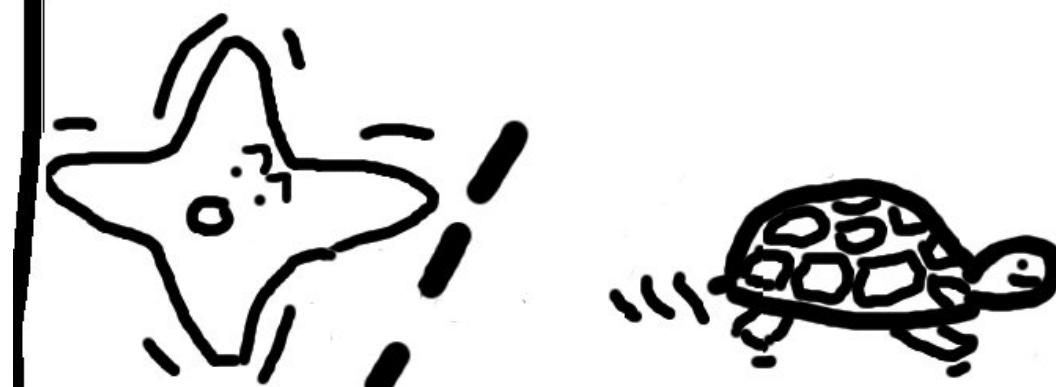
2007 - Jan



Weierstrass has made some progress -
finally Edwards wakes up.

Weierstrass vs. Edwards IV

2007-Jan



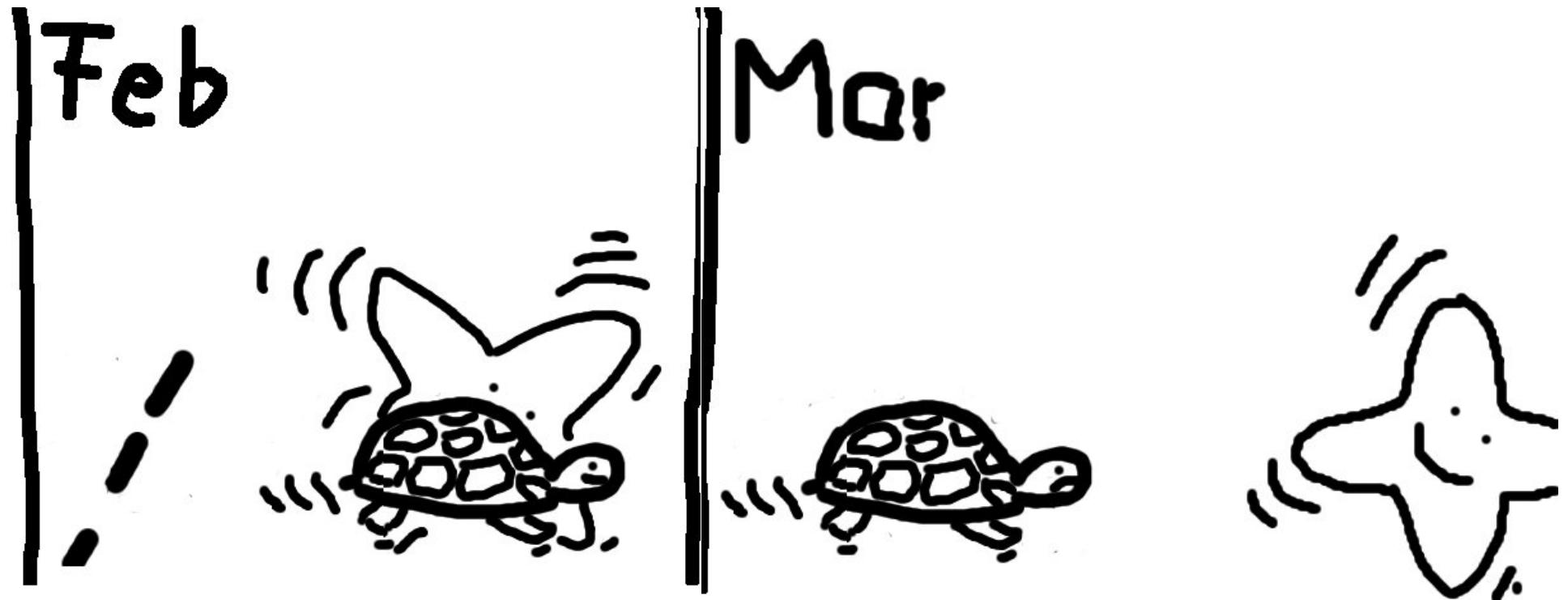
Weierstrass has made some progress –
finally Edwards wakes up.

Feb



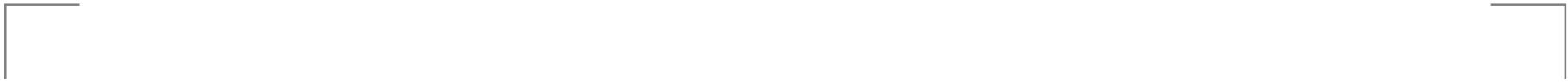
Exciting progress: Edwards
about to overtake!!

Weierstrass vs. Edwards V



Exciting progress: Edwards
about to overtake!!

And the winner is: Edwards!

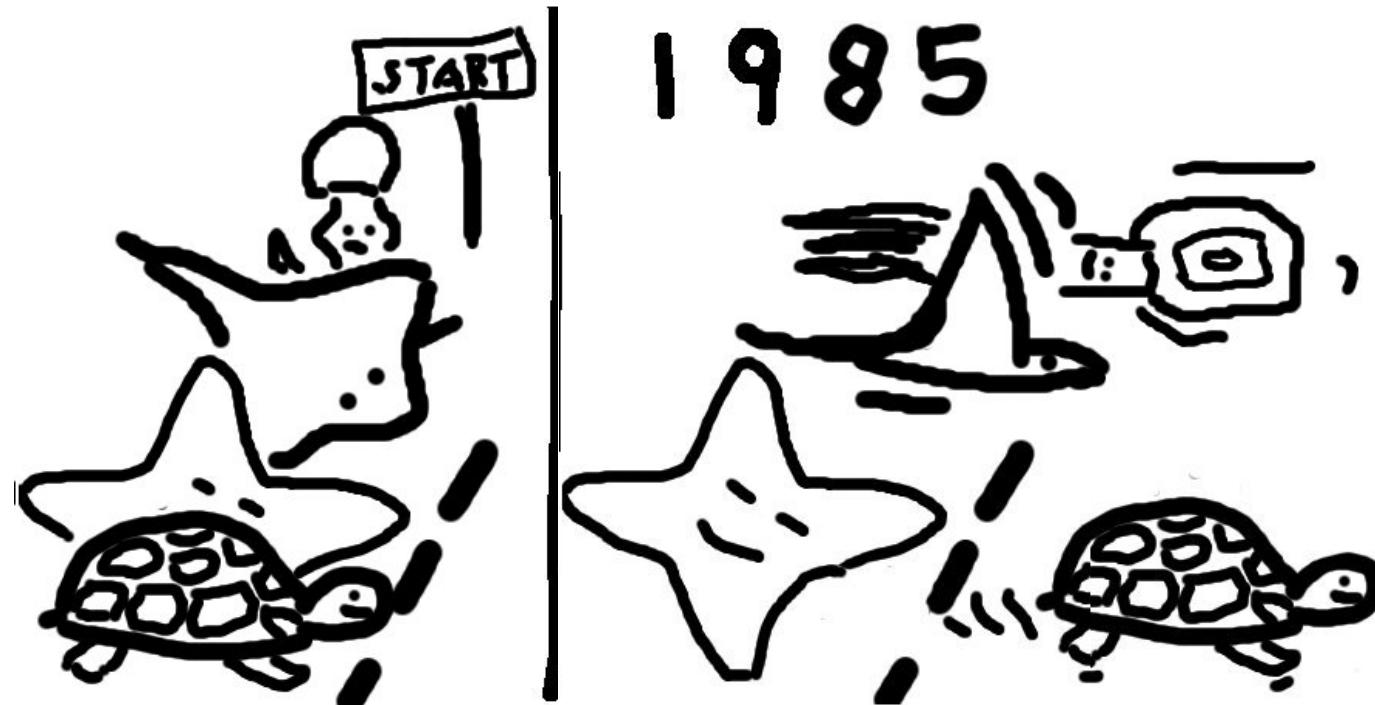


all competitors . . .

All competitors I



All competitors II



All competitors III

1985



2007-Jan

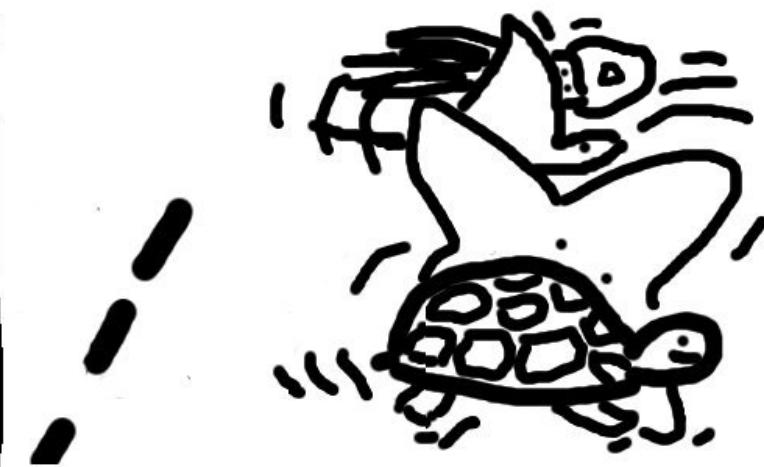


All competitors IV

2007-Jan

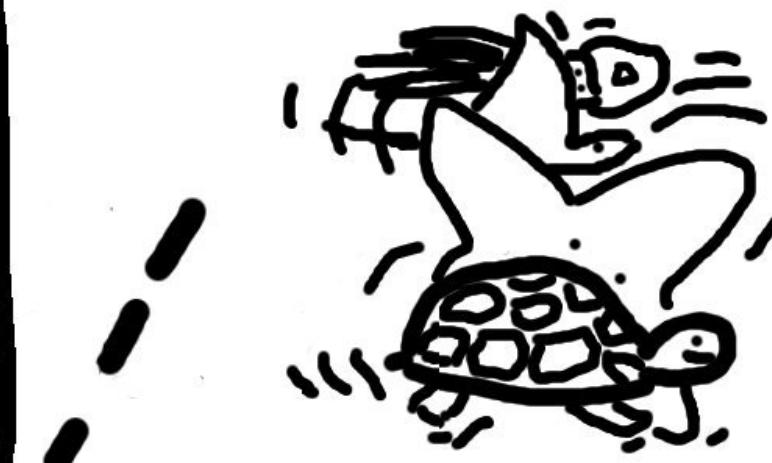


Feb



All competitors V

Feb



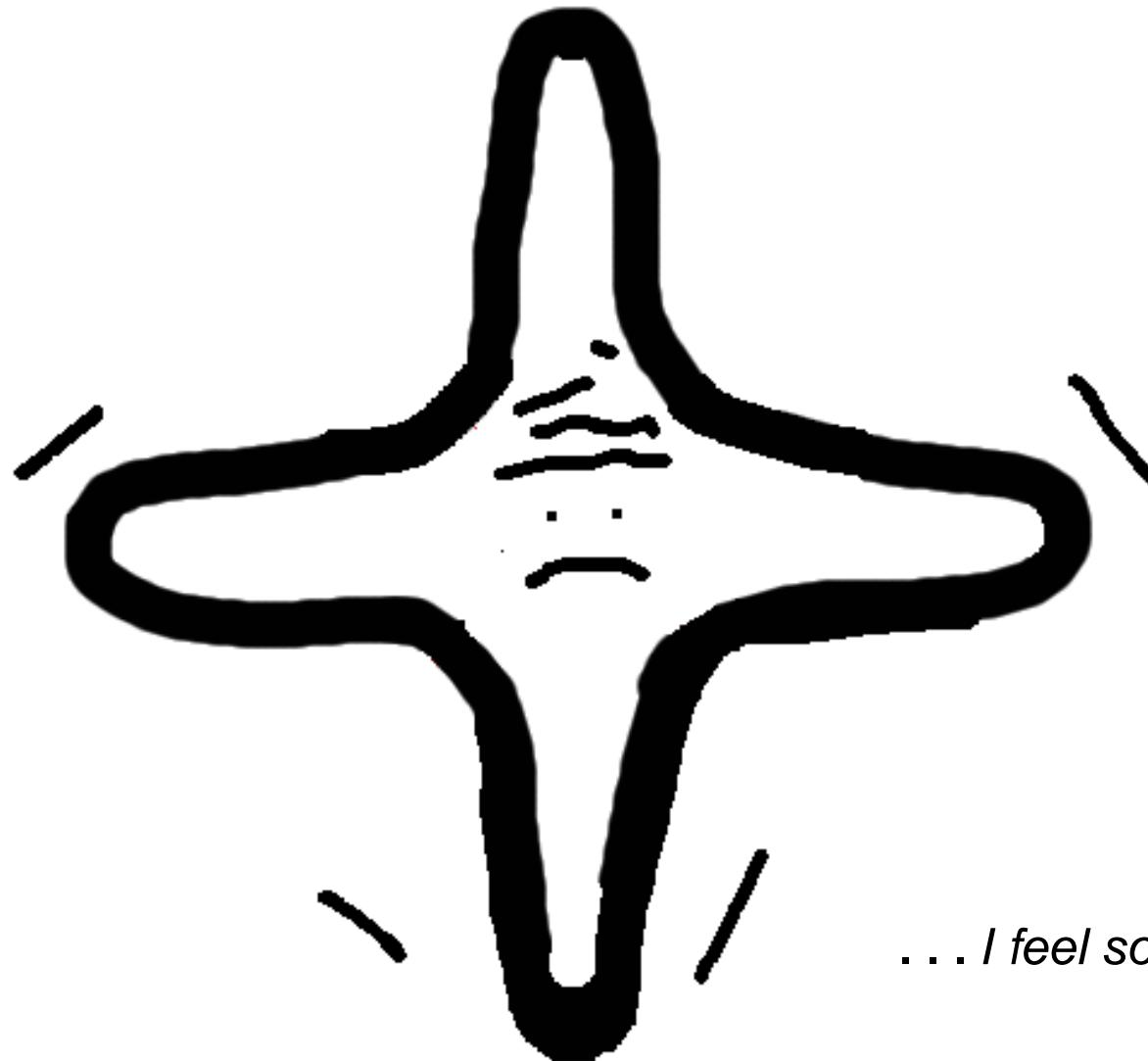
Mar



Read the full story at:

hyperelliptic.org/EFD

One year passes ...



... I feel so odd ...

Exceptions, $2 \neq 0 \dots$

Fix a field k of characteristic different from 2. Fix $c, d \in k$ such that $c \neq 0$, $d \neq 0$, and $dc^4 \neq 1$. Consider the *Edwards addition law*

$$(x_1, y_1), (x_2, y_2) \mapsto \left(\frac{x_1 y_2 + y_1 x_2}{c(1 + dx_1 x_2 y_1 y_2)}, \frac{y_1 y_2 - x_1 x_2}{c(1 - dx_1 x_2 y_1 y_2)} \right)$$

$$x^2 + y^2 = a^2(1 + x^2 y^2), a^5 \neq a$$

describes an elliptic curve over
field k of odd characteristic.

Theorem 2.1. Let k be a field in which $2 \neq 0$. Let E be an elliptic curve over k such that the group $E(k)$ has an element of order 4. Then

How can there be an incomplete set of complete curves???

After extensive (finite) field studies ...

(joint work with Reza Rezaeian Farashahi)

- Assume $d_1, d_2 \in \mathbb{F}_{2^n}$, $d_1 \neq 0$, $\text{Tr}(d_2) = 1$. Then

$$d_1(x + y) + d_2(x + y)^2 = xy + xy(x + y) + x^2y^2$$

describes an elliptic curve.

(Curve shape is symmetric and has highest term x^2y^2 like 'classic' Edwards curves $x^2 + y^2 = 1 + dx^2y^2$.)

- Neutral element is $(0, 0)$. Negative of (x, y) is (y, x) .
- The addition law on this curve is complete! It works for adding arbitrary points – doubling, adding negatives, adding the neutral element, ...
- **Every** ordinary elliptic curve over \mathbb{F}_{2^n} is birationally equivalent to a complete binary Edwards curve.

Timeline (after some early aborts)

- February 13th, 2008: Binary Edwards curves born
- February 15th, 2008: Binary Edwards takes first steps
- February 15th, 2008: Binary Edwards curves are complete (!) for $d_2 = 1$ and n odd.
- February 16th, 2008: Binary Edwards plays with d_1, d_2
- February 20th, 2008: Binary Edwards adds differentially
- February 29th, 2008: Binary Edwards reaches all ordinary curves
- March 31st, 2008: Intel announces support for binary Edwards curves (PCMULQDQ in Westmere)
- April 16th, 2008: Sun announces support for binary Edwards curves in Rock.

Operation counts

These curves are the first binary curves to offer complete addition laws. They are also surprisingly fast:

- ADD on binary Edwards curves takes $21M+1S+4D$, mADD takes $13M+3S+3D$.
- Latest results ADD in $18M+2S+7D$.
- Differential addition ($P + Q$ given P, Q , and $Q - P$) takes $8M+1S+2D$; mixed version takes $6M+1S+2D$.
- Differential addition+doubling (typical step in Montgomery ladder) takes $8M+4S+2D$; mixed version takes $6M+4S+2D$.

See our preprint (ePrint 2008/171) or

cr.yp.to/papers.html#edwards2

for full details, speedups for $d_1 = d_2$, how to choose small coefficients, affine formulas, ...

Comparison with other doubling formulas

Assume curves are chosen with small coefficients.

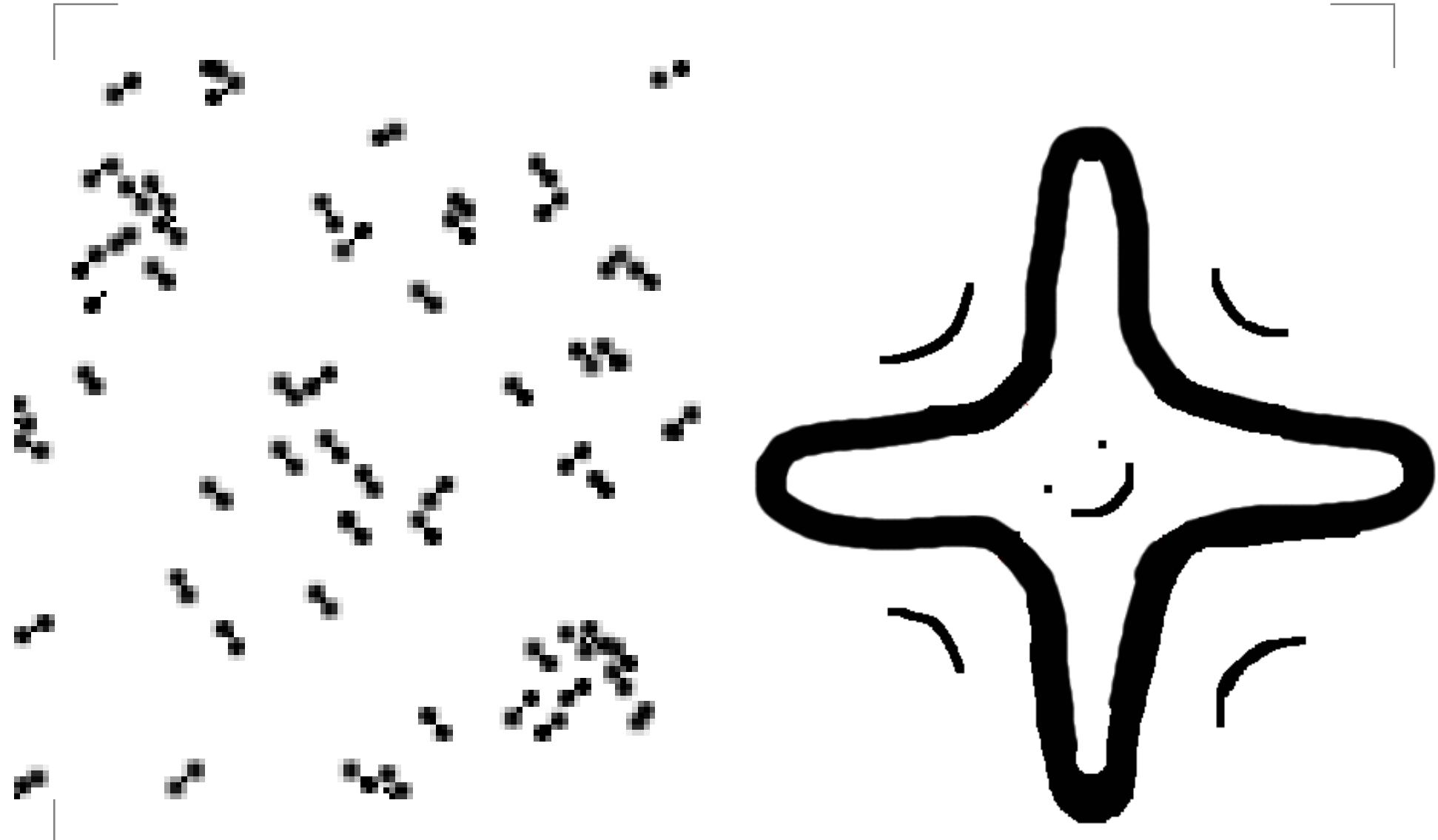
System	Cost of doubling
Projective	7M+4S; see HEHCC
Jacobian	4M+5S; see HEHCC
Lopez-Dahab	3M+5S; Lopez-Dahab
Edwards	2M+6S; new, complete
Lopez-Dahab $a_2 = 1$	2M+5S; Kim-Kim

Explicit-Formulas Database

www.hyperelliptic.org/EFD

for characteristic 2 is in preparation; our paper already has some speed-ups for Lopez-Dahab coordinates.

Happy End!



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