Edwards Coordinates for Elliptic Curves, part 1

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joint work with Daniel J. Bernstein

10.11.2007

Do you know how to add on a circle?

Let k be a field with $2 \neq 0$.

$$\{(x, y) \in k \times k | x^2 + y^2 = 1\}$$

Do you know how to add on a circle?

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is a commutative group with $(x_1, y_1) \oplus (x_2, y_2) = (x_3, y_3)$, where

 $x_3 = x_1y_2 + y_1x_2$ and $y_3 = y_1y_2 - x_1x_2$.

- Polar coordinates and trigonometric identities readily show that the result is on the curve.
- Associativity of the addition boils down to associativity of addition of angles.
- Look, an addition law!
- But it's not elliptic; index calculus work efficiently.

An elliptic curve:

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$$x^2 + y^2 = a^2(1 + x^2y^2)$$

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elliptic? use $z = y(1 - a^2x^2)/a$ to obtain

$$z^{2} = x^{4} - (a^{2} + 1/a^{2})x^{2} + 1.$$

Let k be a field with $2 \neq 0$ and let $a \in k$ with $a^5 \neq a$. There is an – almost everywhere defined – operation on the set

$$\{(x,y) \in k \times k | x^2 + y^2 = a^2(1 + x^2y^2)\}$$

as

$$(x_1, y_1) \oplus (x_2, y_2) = (x_3, y_3)$$

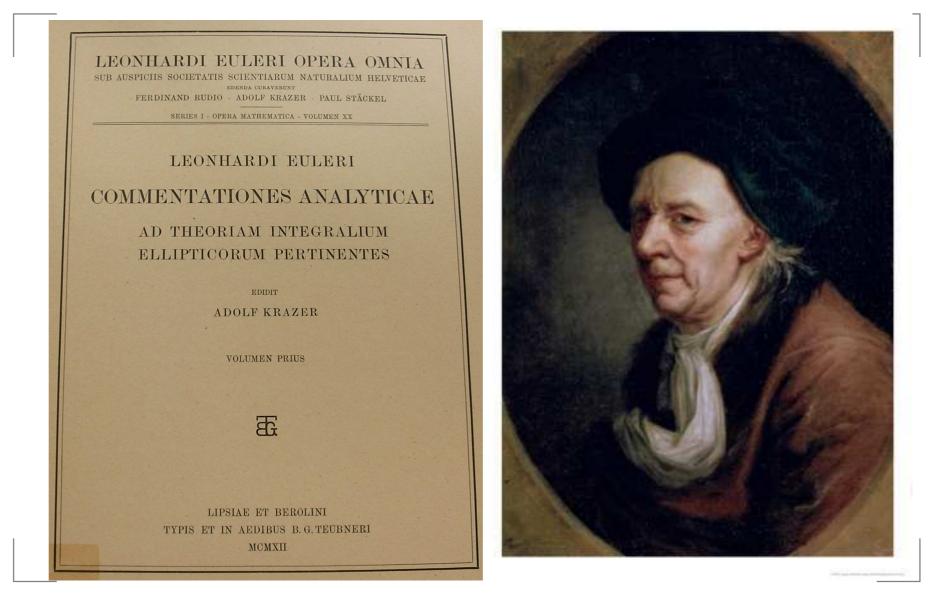
defined by the Edwards addition law

$$x_3 = \frac{x_1y_2 + y_1x_2}{a(1 + x_1x_2y_1y_2)}$$
 and $y_3 = \frac{y_1y_2 - x_1x_2}{a(1 - x_1x_2y_1y_2)}$

Numerators like in addition on circle!

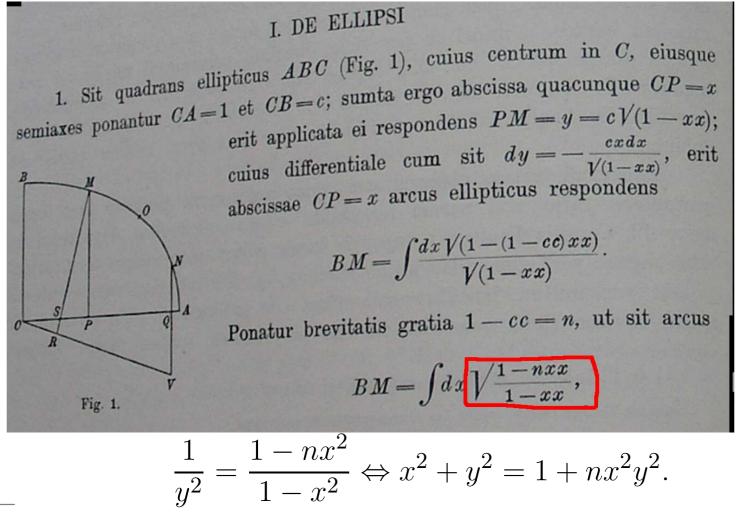
Where do these curves come from?

Long, long ago ...



Euler 1761

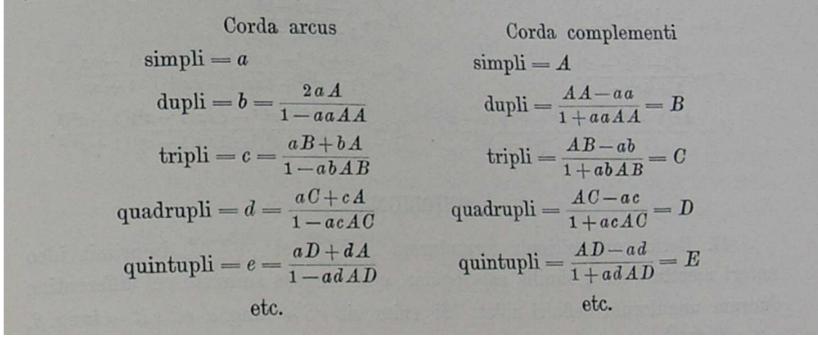
" Observationes de Comparatione Arcuum Curvarum Irrectificabilium"



Euler 1761

COROLLARIUM 3

43. Inventio ergo cordarum arcuum quorumvis multiplorum una cum cordis complementi ita se habebit:



Euler gives doubling and (special) addition for (a, A) on $a^2 + A^2 = 1 - a^2 A^2$.

Gauss, posthumously

ELEGANTIORES INTEGRALIS $\int \frac{dx}{\sqrt{1-x^4}}$ PROPRIETATES. [2.] $1 = ss + cc + sscc \quad sive \quad 2 = (1 + ss)(1 + cc) = (\frac{1}{ss} - 1)(\frac{1}{cc} - 1)(\frac{1}$ $s = \sqrt{\frac{1-cc}{1+cc}}, \quad c = \sqrt{\frac{1-ss}{1+ss}}$ $\sin \operatorname{lemn} \left(a \pm b \right) = \frac{s \, c' \pm s' c}{1 \pm s \, c \, s' \, c'}$ $\cos \operatorname{lemn}(a \pm b) = \frac{cc' \mp ss'}{1 + ss'cc'}$ sin lemn(-a) = -sin lemn a, cos lemn(-a) = cos lemn a $\sin \operatorname{lemn} k \varpi = 0$ $\sin \operatorname{lemn} (k + \frac{1}{2}) \varpi = \pm 1$ $\cos \operatorname{lemn} k \varpi = \pm 1$ $\cos \operatorname{lemn} (k + \frac{1}{2}) \varpi = 0$

Gauss gives general addition for arbitrary points on

 $1 = s^2 + c^2 + s^2 c^2.$

Ex uno plura

- Harold M. Edwards, Bulletin of the AMS, 44, 393–422, 2007 $x^2 + y^2 = a^2(1 + x^2y^2), a^5 \neq a$ describes an elliptic curve.
- Every elliptic curve can be written in this form – over some extension field.
- Ur-elliptic curve $x^2 + y^2 = 1 - x^2 y^2$

needs $\sqrt{-1} \in k$ transform.



Edwards gives above-mentioned addition law for this generalized form, shows equivalence with Weierstrass form, proves addition law, gives theta parameterization ...

Edwards curves over finite fields

- We do not necessarily have $\sqrt{-1} \in k!$ The example curve $x^2 + y^2 = 1 x^2y^2$ from Euler and Gauss is not always an Edwards curve.
- Solution: change the definition of Edwards curves.
- Introduce further parameter d to cover more curves

$$x^{2} + y^{2} = c^{2}(1 + dx^{2}y^{2}), \ c, d \neq 0, dc^{4} \neq 1.$$

- At least one of c, d small: if c⁴d = c⁴d then x² + y² = c²(1 + dx²y²) and x² + y² = c²(1 + dx²y²) isomorphic. We can always choose c = 1 (and do so in the sequel).
- $\bar{c}^4 \bar{d} = (c^4 d)^{-1}$ gives quadratic twist (might be isomorphic).

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

Neutral element is

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• Neutral element is (0, 1), this is an affine point!

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$$-(x_1, y_1) = (-x_1, y_1).$$

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- Neutral element is (0, 1), this is an affine point!
- $-(x_1, y_1) = (-x_1, y_1).$
- (0,-1) has order 2, (±1,0) have order 4,
 so not every elliptic curve can be transformed to an
 Edwards curve over k but every curve with a point of order 4 can!
- Our Asiacrypt 2007 paper makes explicit the birational equivalence between a curve in Edwards form and in Weierstrass form.

See also our newelliptic page.

$$P \oplus Q = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

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- Addition law also works for doubling (compare that to curves in Weierstrass form!)
- Can show: denominator never 0 for non-square d.

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Can show: denominator never 0 for non-square d.
Explicit formulas for addition/doubling:

$$A = Z_1 \cdot Z_2; \ B = A^2; \ C = X_1 \cdot X_2; \ D = Y_1 \cdot Y_2;$$

$$E = (X_1 + Y_1) \cdot (X_2 + Y_2) - C - D; \ F = d \cdot C \cdot D;$$

$$X_{P \oplus Q} = A \cdot E \cdot (B - F); \ Y_{P \oplus Q} = A \cdot (D - C) \cdot (B + F);$$

$$Z_{P \oplus Q} = (B - F) \cdot (B + F).$$

•
$$P \oplus Q = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

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$$Z_{P \oplus Q} = (B - F) \cdot (B + F).$$

Needs 10M + 1S + 1D + 7A.

Strongly unified group operations

- Addition formulas work also for doubling.
- Addition in Weierstrass form $y^2 = x^3 + a_4x + a_6$, involves computation

$$\lambda = \begin{cases} (y_2 - y_1)/(x_2 - x_1) & \text{if } x_1 \neq x_2, \\ (3x_1^2 + a_4)/(2y_1) & \text{else.} \end{cases}$$

division by zero if first form is accidentally used for doubling.

- Strongly unified addition laws remove some checks from the code.
- Help against simple side-channel attacks. Attacker sees uniform sequence of identical group operations, no information on secret scalar given (assuming the field operations are handled appropriately).

Unified Projective coordinates

Brier, Joye 2002

Idea: unify how the slope is computed.

improved in Brier, Déchène, and Joye 2004

$$\lambda = \frac{(x_1 + x_2)^2 - x_1 x_2 + a_4 + y_1 - y_2}{y_1 + y_2 + x_1 - x_2}$$
$$= \begin{cases} \frac{y_1 - y_2}{x_1 - x_2} & (x_1, y_1) \neq \pm (x_2, y_2) \\ \frac{3x_1^2 + a_4}{2y_1} & (x_1, y_1) = (x_2, y_2) \end{cases}$$

Multiply numerator & denominator by $x_1 - x_2$ to see this.

- Proposed formulae can be generalized to projective coordinates.
- Some special cases may occur, but with very low probability, e.g. $x_2 = y_1 + y_2 + x_1$. Alternative equation for this case.

Jacobi intersections

- Chudnovsky and Chudnovsky 1986; Liardet and Smart CHES 2001
- Elliptic curve given as intersection of two quadratics

$$s^2 + c^2 = 1$$
 and $as^2 + d^2 = 1$.

- Points (S : C : D : Z) with (s, c, d) = (S/Z, C/Z, D/Z).
- Neutral element is (0, 1, 1).

$$S_3 = (Z_1C_2 + D_1S_2)(C_1Z_2 + S_1D_2) - Z_1C_2C_1Z_2 - D_1S_2S_1D_2$$

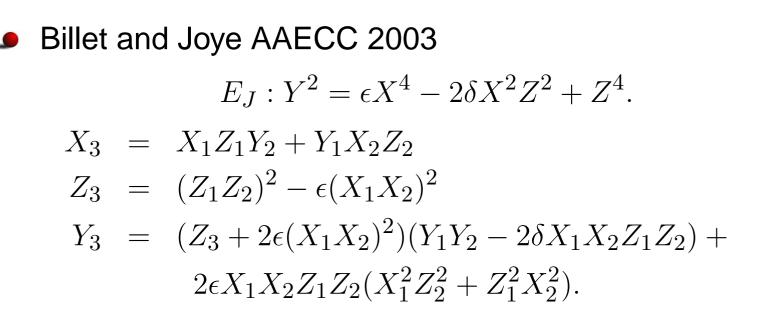
$$C_3 = Z_1 C_2 C_1 Z_2 - D_1 S_2 S_1 D_2$$

 $D_3 = Z_1 D_1 Z_2 D_2 - a S_1 C_1 S_2 C_2$

$$Z_3 = Z_1 C_2^2 + D_1 S_2^2.$$

Unified formulas need 13M + 2S + 1D.

Jacobi quartics



- Unified formulas need 10M+3S+D+2E
- Can have ϵ or δ small
- Needs point of order 2; for $\epsilon = 1$ the group order is divisible by 4.
- Some recent speed ups due to Duquesne and to Hisil, Carter, and Dawson.

Hessian curves

$$E_H: X^3 + Y^3 + Z^3 = cXYZ.$$

 $\begin{array}{ll} \text{Addition: } P \neq \pm Q & \text{Doubling } P = Q \neq -P \\ X_3 = X_2 Y_1^2 Z_2 - X_1 Y_2^2 Z_1 & X_3 = Y_1 (X_1^3 - Z_1^3) \\ Y_3 = X_1^2 Y_2 Z_2 - X_2^2 Y_1 Z_1 & Y_3 = X_1 (Z_1^3 - Y_1^3) \\ Z_3 = X_2 Y_2 Z_1^2 - X_1 Y_1 Z_2^2 & Z_3 = Z_1 (Y_1^3 - X_1^3) \end{array}$

- Curves were first suggested for speed
- Joye and Quisquater show

 $[2](X_1:Y_1:Z_1) = (Z_1:X_1:Y_1) \oplus (Y_1:Z_1:X_1)$

- Unified formulas need 12M.
- Doubling is done by an addition, but not automatically only unified, not strongly unified.

Unified addition law

- Unified formulas introduced as countermeasure against side-channel attacks but useful in general.
- Strongly unified addition laws indeed remove the check for $P \neq Q$ before addition.
- Some systems allow to omit the check $P \neq -Q$ before addition.
- Most systems still have exceptional cases.
- No surprise:

"The smallest cardinality of a complete system of addition laws on *E* equals two." (Theorem 1 in Wieb Bosma, Hendrik W. Lenstra, Jr., J. Number Theory **53**, 229–240, 1995)

Bosma/Lenstra give such system; similar to unified projective coordinates.

Complete addition law

- If *d* is not a square then Edwards addition law is complete: For $x_i^2 + y_i^2 = 1 + dx_i^2 y_i^2$, i = 1, 2, always $dx_1x_2y_1y_2 \neq \pm 1$. Outline of proof: If $(dx_1x_2y_1y_2)^2 = 1$ then $(x_1 + dx_1x_2y_1y_2y_1)^2 =$ $dx_1^2y_1^2(x_2 + y_2)^2$. Conclude that *d* is a square. But $d \neq \Box$.
- Edwards addition law allows omitting all checks
 - Neutral element is affine point on curve.
 - Addition works to add P and P.
 - Addition works to add P and -P.
 - Addition just works to add P and any Q.
- Only complete addition law in the literature.
- Bosma/Lenstra strikes over quadratic extension. "Pointless exceptional divisor!"

Fastest unified addition-or-doubling formula

System	Cost of unified addition-or-doubling
Projective	11M+6S+1D; see Brier/Joye '03
Projective if $a_4 = -1$	13M+3S; see Brier/Joye '02
Jacobi intersection	13M+2S+1D; see Liardet/Smart '01
Jacobi quartic ($\epsilon = 1$)	10M+3S+1D; see Billet/Joye '01
Hessian	12M; see Joye/Quisquater '01
Edwards	10M+1S+1D

- Exactly the same formulae for doubling (no re-arrangement like in Hessian; no if-else)
- No exceptional cases if d is not a square.
- Operation counts as in Asiacrypt'07 paper.
- See EFD hyperelliptic.org/EFD.

What if we know that we double?

$$[2]P = \left(\frac{x_1y_1 + y_1x_1}{1 + dx_1x_1y_1y_1}, \frac{y_1y_1 - x_1x_1}{1 - dx_1x_1y_1y_1}\right)$$
$$= \left(\frac{2x_1y_1}{1 + d(x_1y_1)^2}, \frac{y_1^2 - x_1^2}{1 - d(x_1y_1)^2}\right)$$

$$[2]P = \left(\frac{x_1y_1 + y_1x_1}{1 + dx_1x_1y_1y_1}, \frac{y_1y_1 - x_1x_1}{1 - dx_1x_1y_1y_1}\right)$$
$$= \left(\frac{2x_1y_1}{1 + d(x_1y_1)^2}, \frac{y_1^2 - x_1^2}{1 - d(x_1y_1)^2}\right)$$

Use curve equation $x^2 + y^2 = 1 + dx^2y^2$.

$$[2]P = \left(\frac{x_1y_1 + y_1x_1}{1 + dx_1x_1y_1y_1}, \frac{y_1y_1 - x_1x_1}{1 - dx_1x_1y_1y_1}\right)$$
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$$= \left(\frac{2x_1y_1}{x_1^2 + y_1^2}, \frac{y_1^2 - x_1^2}{2 - (x_1^2 + y_1^2)}\right)$$

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$$= \left(\frac{2x_1y_1}{x_1^2 + y_1^2}, \frac{y_1^2 - x_1^2}{2 - (x_1^2 + y_1^2)}\right)$$

 $B = (X_1 + Y_1)^2; \ C = X_1^2; \ D = Y_1^2; \ E = C + D; \ H = (c \cdot Z_1)^2;$ $J = E - 2H; \ X_3 = c \cdot (B - E) \cdot J; \ Y_3 = c \cdot E \cdot (C - D); \ Z_3 = E \cdot .$

Inversion-free version needs 3M + 4S + 6A.

Very fast doubling formulae

System	Cost of doubling
Projective	5M+6S+1D; EFD
Projective if $a_4 = -3$	7M+3S; EFD
Hessian	7M+1S; see Hisil/Carter/Dawson '07
Doche/Icart/Kohel-3	2M+7S+2D; see Doche/Icart/Kohel '06
Jacobian	1M+8S+1D; EFD
Jacobian if $a_4 = -3$	3M+5S; see DJB '01
Jacobi quartic	2M+6S+2D; see Hisil/Carter/Dawson '07
Jacobi intersection	3M+4S; see Liardet/Smart '01
Edwards	3M+4S;
Doche/Icart/Kohel-2	2M+5S+2D; see Doche/Icart/Kohel '06

- Edwards fastest for general curves, no D.
- Operation counts as in our Asiacrypt paper.

Fastest addition formulae

System	Cost of addition
Doche/Icart/Kohel-2	12M+5S+1D; see Doche/Icart/Kohel '06
Doche/Icart/Kohel-3	11M+6S+1D; see Doche/Icart/Kohel '06
Jacobian	11M+5S; EFD
Jacobi intersection	13M+2S+1D; see Liardet/Smart '01
Projective	12M+2S; HECC
Jacobi quartic	10M+3S+1D; see Billet/Joye '03
Hessian	12M; see Joye/Quisquater '01
Edwards	10M+1S+1D

- EFD and full paper also contain costs for mixed addition (mADD) and re-additions (reADD).
- reADD: non-mixed ADD where one point has been added before and computations have been cached.

Single-scalar multiplication using NAF

System	1 DBL, 1/3 mADD			
Projective	8M+6.67S+1D			
Projective if $a_4 = -3$	10M+3.67S			
Hessian	10.3M+1S			
Doche/Icart/Kohel-3	4.33M+8.33S+2.33D			
Jacobian	3.33M+9.33S+1D			
Jacobian if $a_4 = -3$	5.33M+6.33S			
Jacobi intersection	6.67M+4.67S+0.333D			
Jacobi quartic	4.67M+7S+2.33D			
Doche/Icart/Kohel-2	4.67M+6.33S+2.33D			
Edwards	6M+4.33S+0.333D			
For comparison. Montgomery arithmetic takes 5M+4S				

For comparison: Montgomery arithmetic takes 5M+4S+1D per bit.

Signed width-4 sliding windows

These counts include the precomputations.

System	0.98 DBL, 0.17 reADD, 0.025 mADD, 0.0035 A		
Projective	7.17M+6.28S+0.982D		
Projective if $a_4 = -3$	9.13M+3.34S		
Doche/Icart/Kohel-3	3.84M+7.99S+2.16D		
Hessian	9.16M+0.982S		
Jacobian	2.85M+8.64S+0.982D		
Jacobian if $a_4 = -3$	4.82M+5.69S		
Doche/Icart/Kohel-2	4.2M+5.86S+2.16D		
Jacobi quartic	3.69M+6.48S+2.16D		
Jacobi intersection	5.09M+4.32S+0.194D		
Edwards	4.86M+4.12S+0.194D		
Montgomery takes 5M+4S+1D per bit. Edwards solidly faster!			

Inverted Edwards coordinates

- Latest news (Bernstein/Lange, to appear at AAECC 2007): inverted Edwards coordinates are even faster strongly unified system – but not complete.
- Using the representation $(X_1 : Y_1 : Z_1)$ for the affine point $(Z_1/X_1, Z_1/Y_1)$ $(X_1Y_1Z_1 \neq 0)$ gives operation counts:
 - Doubling takes 3M + 4S + 1D.
 - Addition takes 9M + 1S + 1D.
- This saves 1M for each addition compared to standard Edwards coordinates.
- New speed leader: inverted Edwards coordinates.

Different coordinate systems

For coordinate systems we could find, the group law, operation counts (and improvements) for the explicit formulas, MAGMA-based proofs (sorry, not SAGE) of their correctness, lots of entertainment visit the

Explicit Formulas Database

http://www.hyperelliptic.org/EFD

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

What if denominators are 0?

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

What if denominators are 0?

Answer: They are never 0 if d is not a square in k.

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

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Intuitive explanation:

The points (1:0:0) and (0:1:0) are singular. They correspond to four points on the desingularization of the curve; but those four points are defined over $k(\sqrt{d})$.

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

What if denominators are 0?

Answer: They are never 0 if *d* is not a square in *k*. Explicit proof: Let $(x_1, y_1), (x_2, y_2)$ be on curve, i.e., if $x_i^2 + y_i^2 = 1 + dx_i^2 y_i^2$. Write $\epsilon = dx_1 x_2 y_1 y_2$ and suppose $\epsilon \in \{-1, 1\}$. Then $x_1, x_2, y_1, y_2 \neq 0$ and $dx_1^2 y_1^2 (x_2^2 + y_2^2) = dx_1^2 y_1^2 + d^2 x_1^2 y_1^2 x_2^2 y_2^2$

$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

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$$(x_1, y_1) \oplus (x_2, y_2) = \left(\frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2}, \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}\right)$$

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