High-speed Diffie-Hellman, part 1

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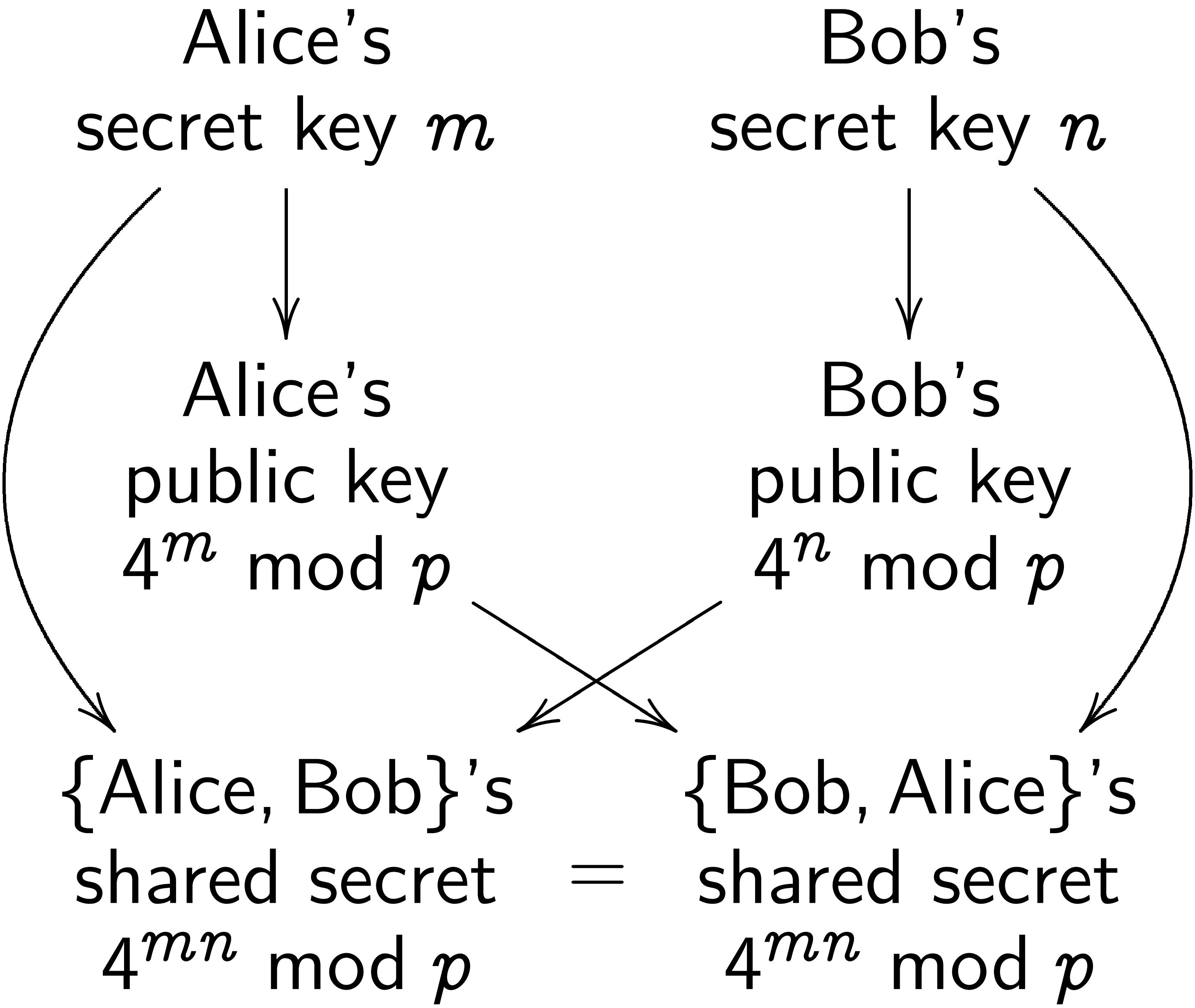
Can quickly compute $4^n \bmod 2^{262} - 5081$ given $n \in \{0, 1, 2, \ldots, 2^{256} - 1\}$.

Similarly, can quickly compute $4^{mn} \bmod 2^{262} - 5081$ given n and $4^m \bmod 2^{262} - 5081$.

"Discrete-logarithm problem": given $4^n \bmod 2^{262} - 5081$, find n. Is this easy to solve?

Diffie-Hellman secret-sharing system using $p = 2^{262} - 5081$:



Can attacker find $4^{mn} \bmod p$?

Bad news: DLP can be solved at surprising speed! Attacker can find m and n by "index calculus."

To protect against this attack, replace $2^{262} - 5081$ with a much larger prime. *Much* slower arithmetic.

Alternative: Elliptic-curve cryptography. Replace $\{1, 2, \ldots, 2^{262} - 5082\}$ with a comparable-size "safe elliptic-curve group." *Somewhat* slower arithmetic.

An elliptic curve over **R**

Consider all pairs
of real numbers x, y
such that $y^2 - 5xy = x^3 - 7$.

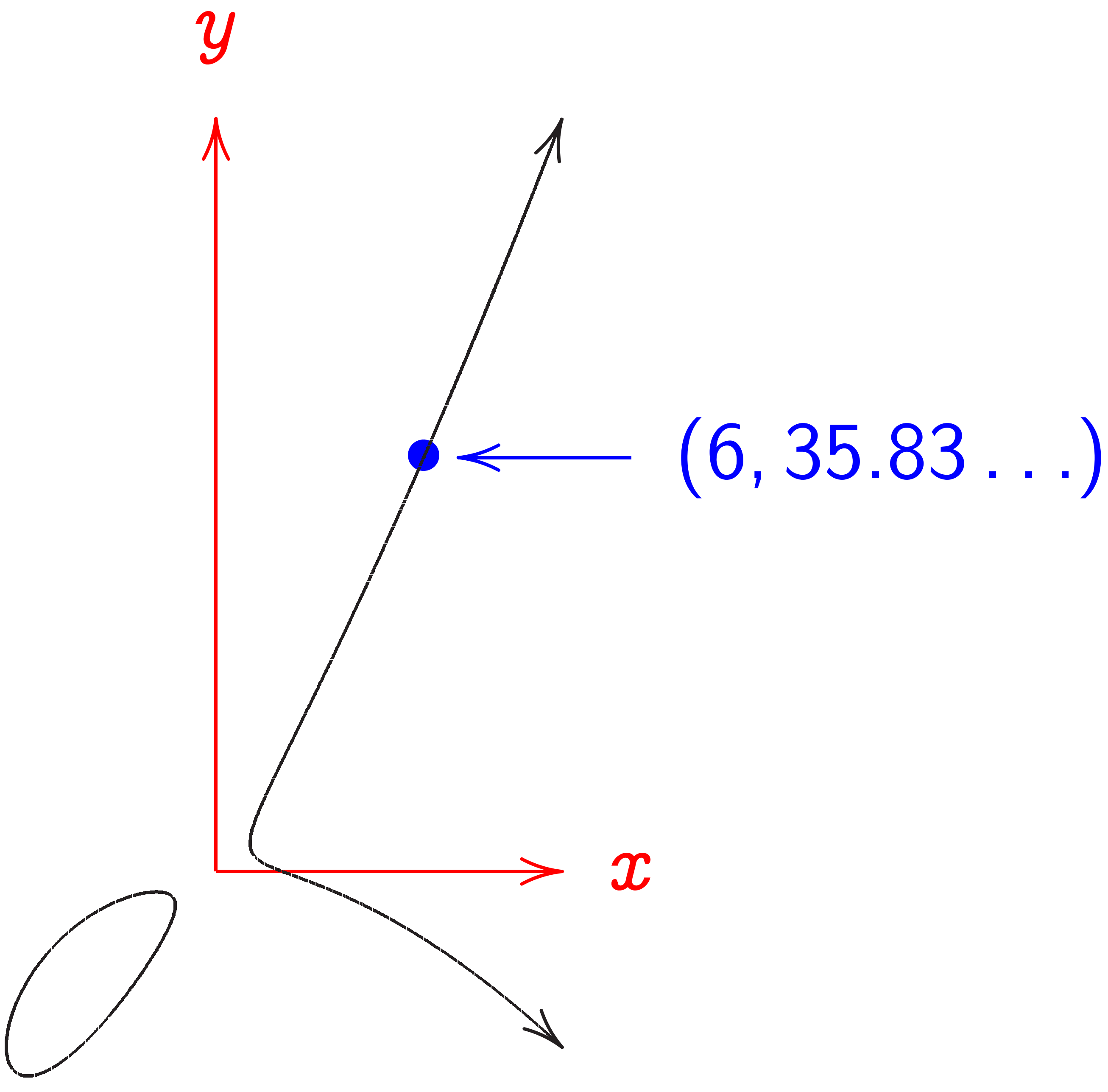
The "points on the elliptic curve
$y^2 - 5xy = x^3 - 7$ over **R**"
are those pairs and
one additional point, ∞.

i.e. The set of points is

$$\{(x, y) \in \mathbf{R} \times \mathbf{R} : \\ y^2 - 5xy = x^3 - 7\} \cup \{\infty\}.$$

(**R** is the set of real numbers.)

Graph of this set of points:



Don't forget ∞.

Visualize ∞ as top of y axis.

There is a standard definition of $0, -, +$ on this set of points.

Magical fact: The set of points is a "commutative group"; i.e., these operations $0, -, +$ satisfy every identity satisfied by $\mathbf{Z}$.

e.g. All $P, Q, R \in \mathbf{Z}$ satisfy $(P + Q) + R = P + (Q + R)$, so all curve points P, Q, R satisfy $(P + Q) + R = P + (Q + R)$.

($\mathbf{Z}$ is the set of integers.)

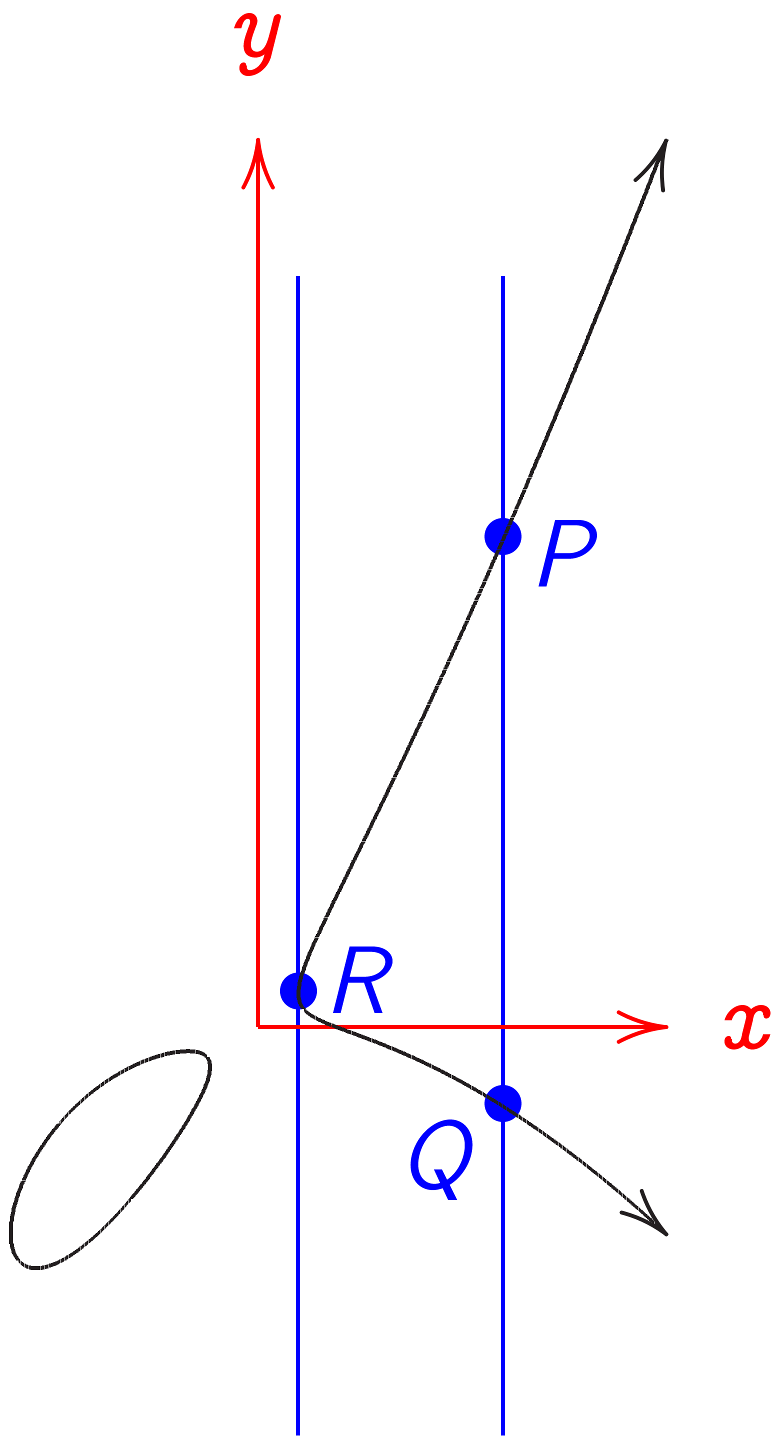
Visualizing the group law

$0 = \infty$; $-\infty = \infty$.

Distinct curve points P, Q
on a vertical line
have $-P = Q$;
$P + Q = 0 = \infty$.

A curve point R
with a vertical tangent line
has $-R = R$;
$R + R = 0 = \infty$.

Here $-P = Q$, $-Q = P$, $-R = R$:

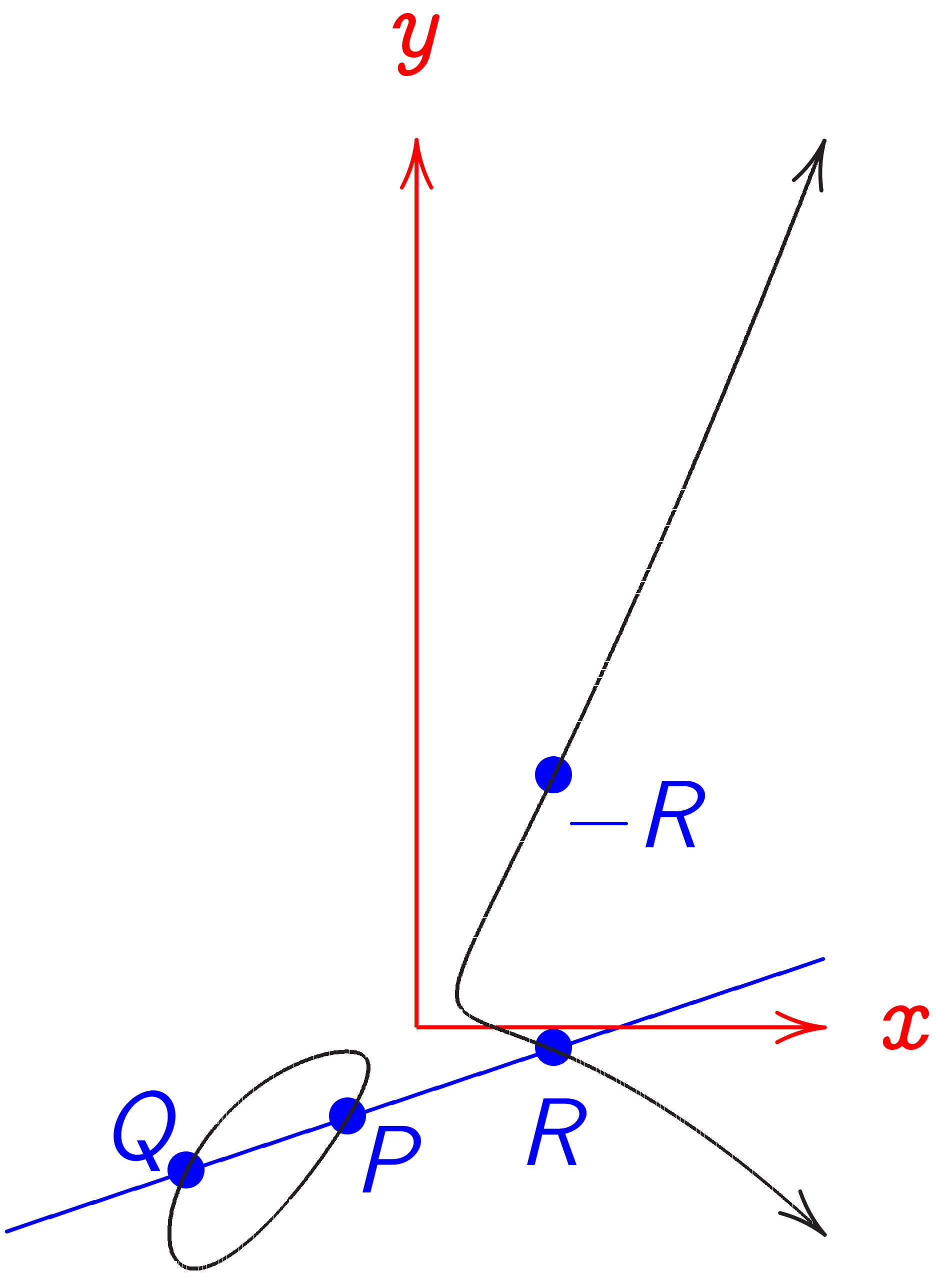


Distinct curve points P, Q, R
on a line
have $P + Q = -R$;
$P + Q + R = 0 = \infty$.

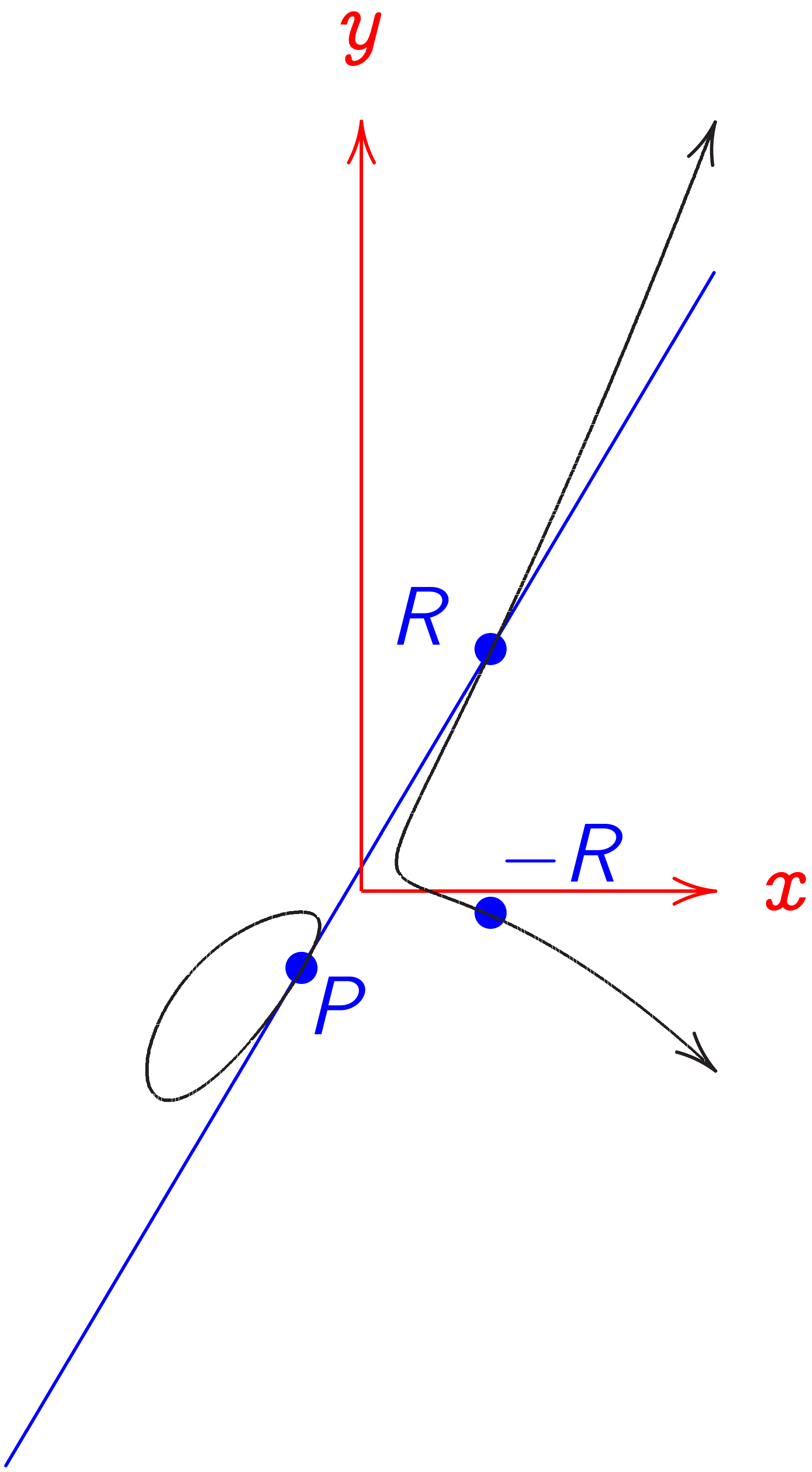
Distinct curve points P, R
on a line tangent at P
have $P + P = -R$;
$P + P + R = 0 = \infty$.

A non-vertical line
with only one curve point P
has $P + P = -P$;
$P + P + P = 0$.

Here $P + Q = -R$:



Here $P + P = -R$:



Curve addition formulas

Easily find formulas for $+$
by finding formulas for lines
and for curve-line intersections.

$x \neq x'$: $(x, y) + (x', y') = (x'', y'')$
where $\lambda = (y' - y)/(x' - x)$,
$x'' = \lambda^2 - 5\lambda - x - x'$,
$y'' = 5x'' - (y + \lambda(x'' - x))$.

$2y \neq 5x$: $(x, y) + (x, y) = (x'', y'')$
where $\lambda = (5y + 3x^2)/(2y - 5x)$,
$x'' = \lambda^2 - 5\lambda - 2x$,
$y'' = 5x'' - (y + \lambda(x'' - x))$.

$(x, y) + (x, 5x - y) = \infty$.

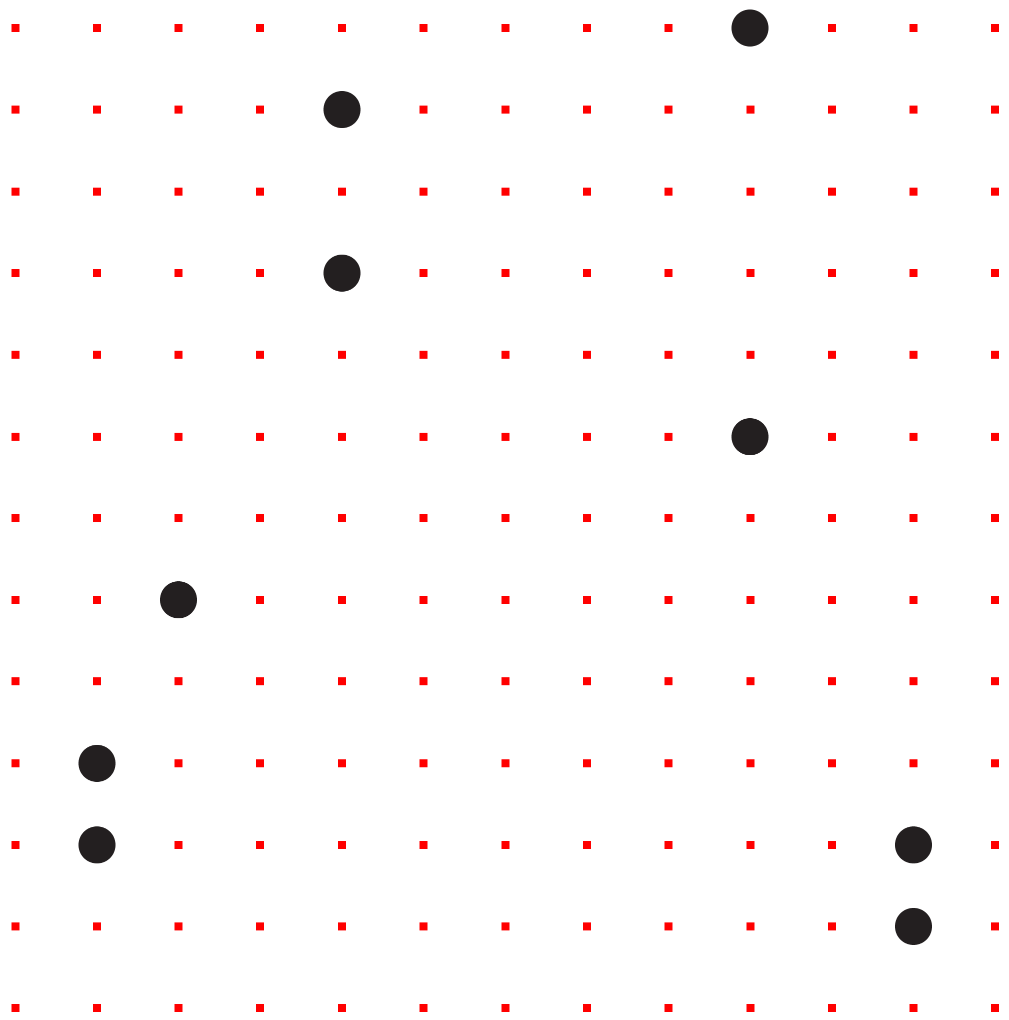
An elliptic curve over $\mathbf{Z}/13$

Consider the prime field
$\mathbf{Z}/13 = \{0, 1, 2, \ldots, 12\}$
with $-, +, \cdot$ defined mod 13.

The "set of points on the elliptic curve $y^2 - 5xy = x^3 - 7$ over $\mathbf{Z}/13$" is

$$\{(x, y) \in \mathbf{Z}/13 \times \mathbf{Z}/13 : \\ y^2 - 5xy = x^3 - 7\} \cup \{\infty\}.$$

Graph of this set of points:



As before, don't forget ∞.

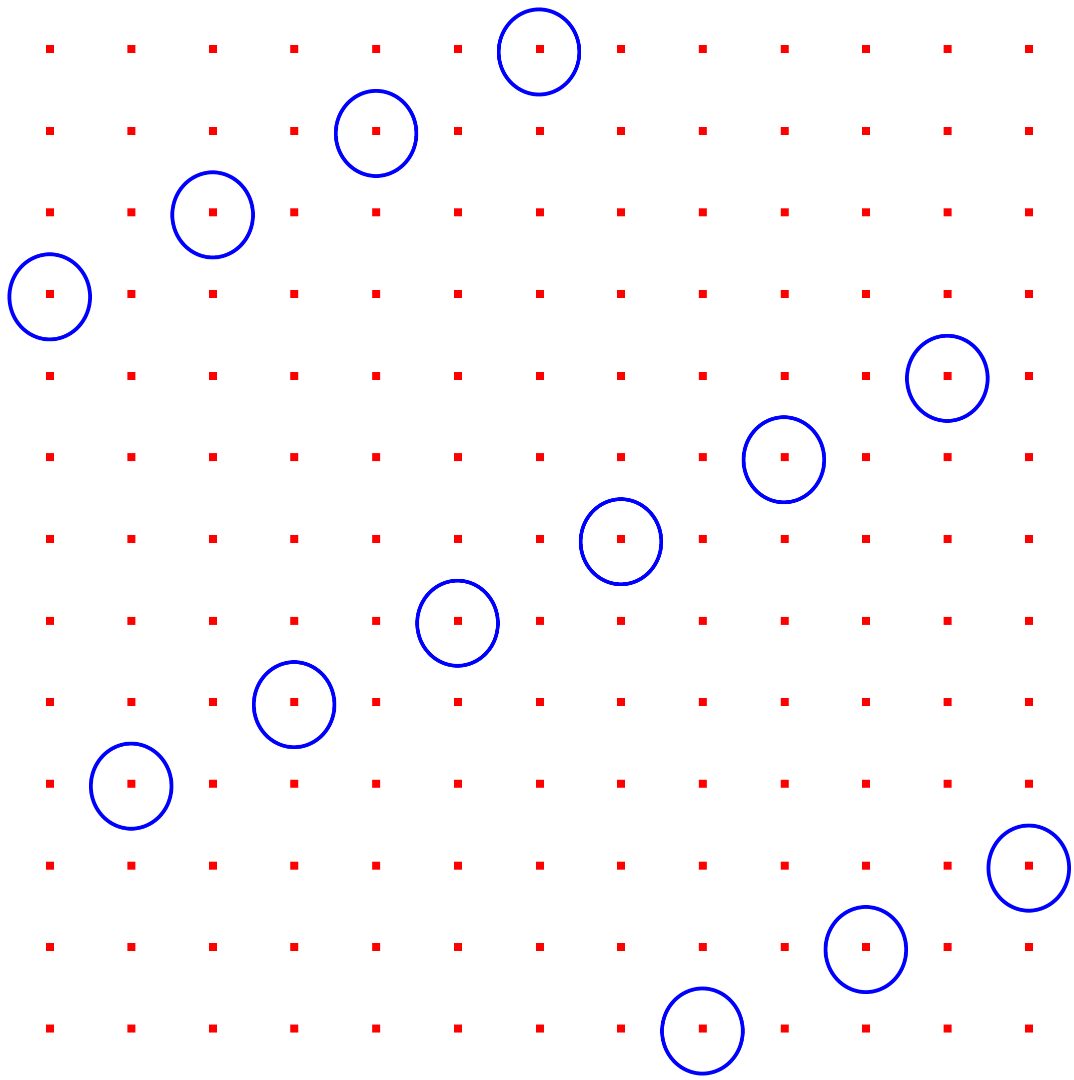
The set of curve points
is a commutative group with
standard definition of $0, -, +$.

Can visualize $0, -, +$ as before.
Replace lines over $\mathbf{R}$
by lines over $\mathbf{Z}/13$.

Warning: tangent is defined by
derivatives; hard to visualize.

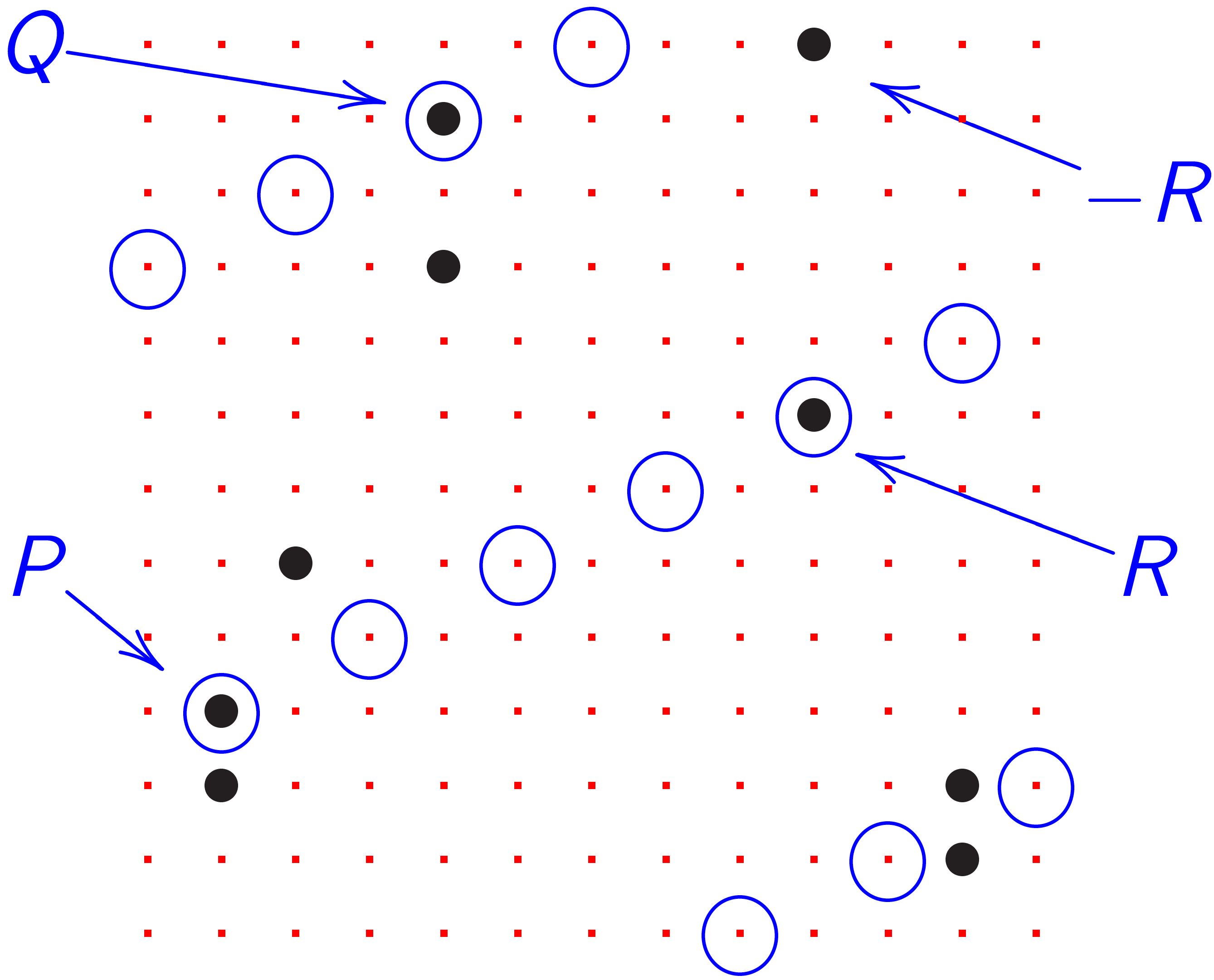
Can define $0, -, +$
using same formulas as before.

Example of line over $\mathbf{Z}/13$:



Formula for this line: $y = 7x + 9$.

$P + Q = -R$:



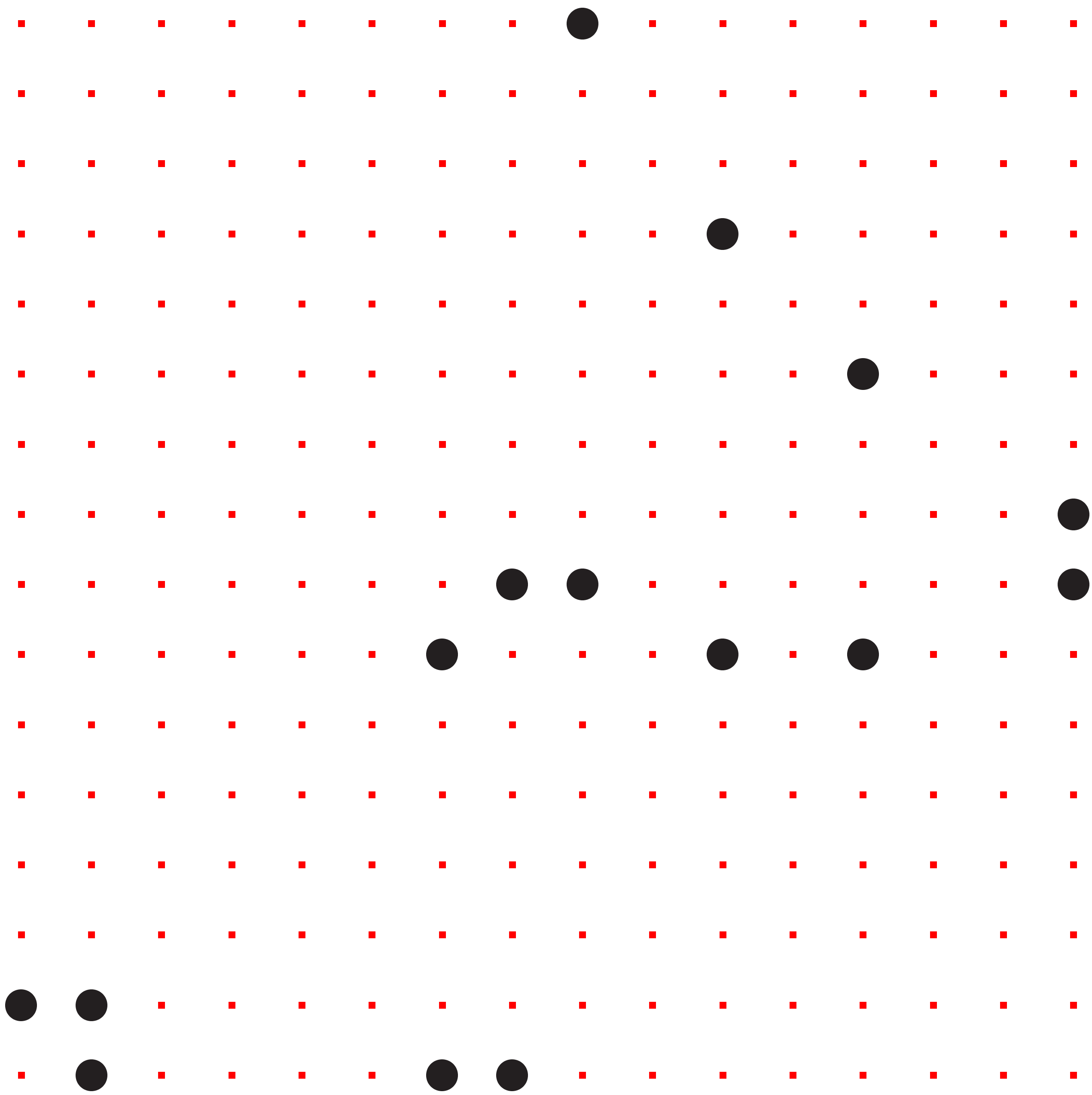
An elliptic curve over $\mathbf{F}_{16}$

Consider the non-prime field

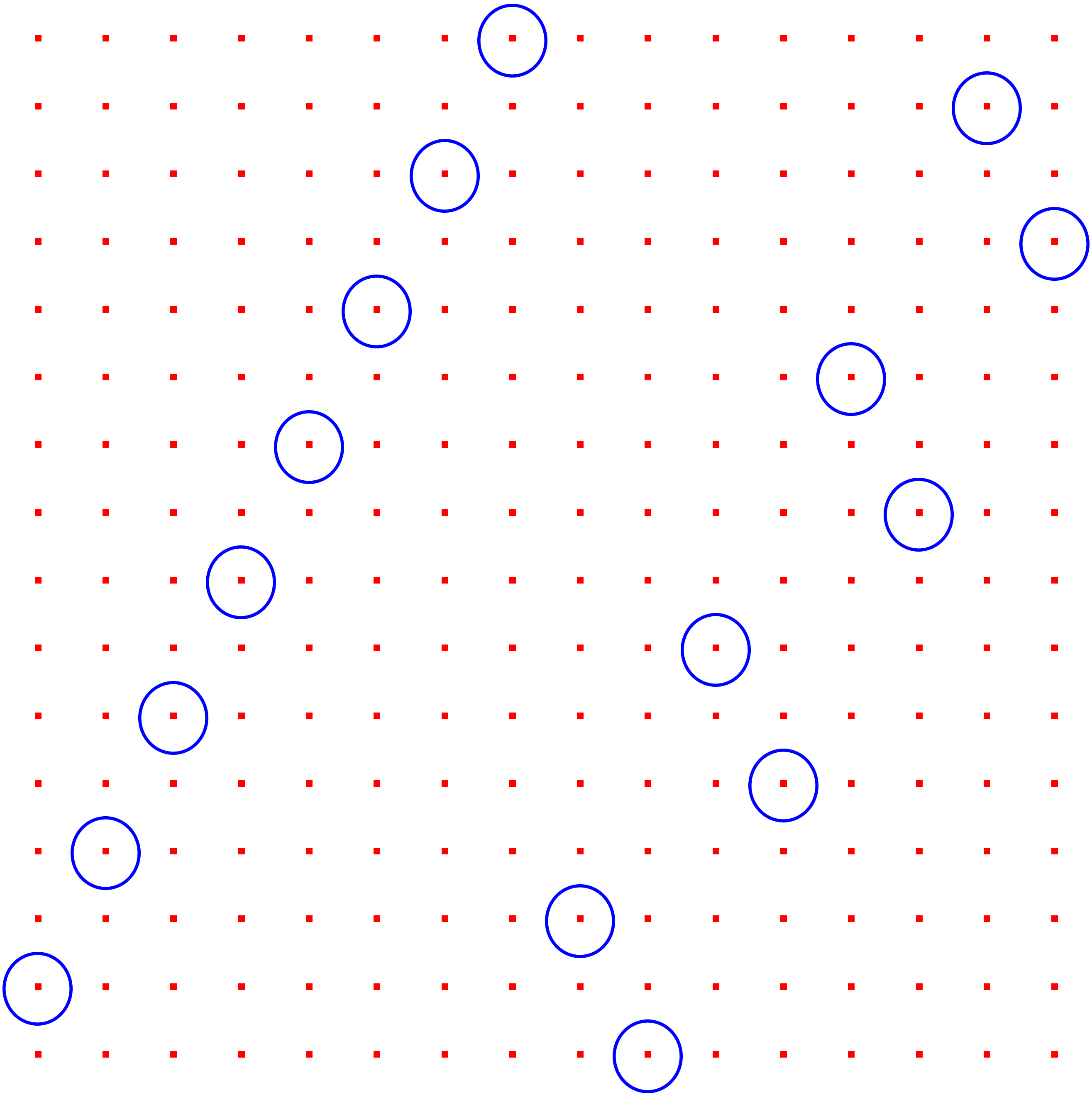
$$(\mathbf{Z}/2)[t]/(t^4 - t - 1) = \{$$
$$0t^3 + 0t^2 + 0t^1 + 0t^0,$$
$$0t^3 + 0t^2 + 0t^1 + 1t^0,$$
$$0t^3 + 0t^2 + 1t^1 + 0t^0,$$
$$0t^3 + 0t^2 + 1t^1 + 1t^0,$$
$$0t^3 + 1t^2 + 0t^1 + 0t^0,$$
$$\vdots$$
$$1t^3 + 1t^2 + 1t^1 + 1t^0\}$$

of size $2^4 = 16$.

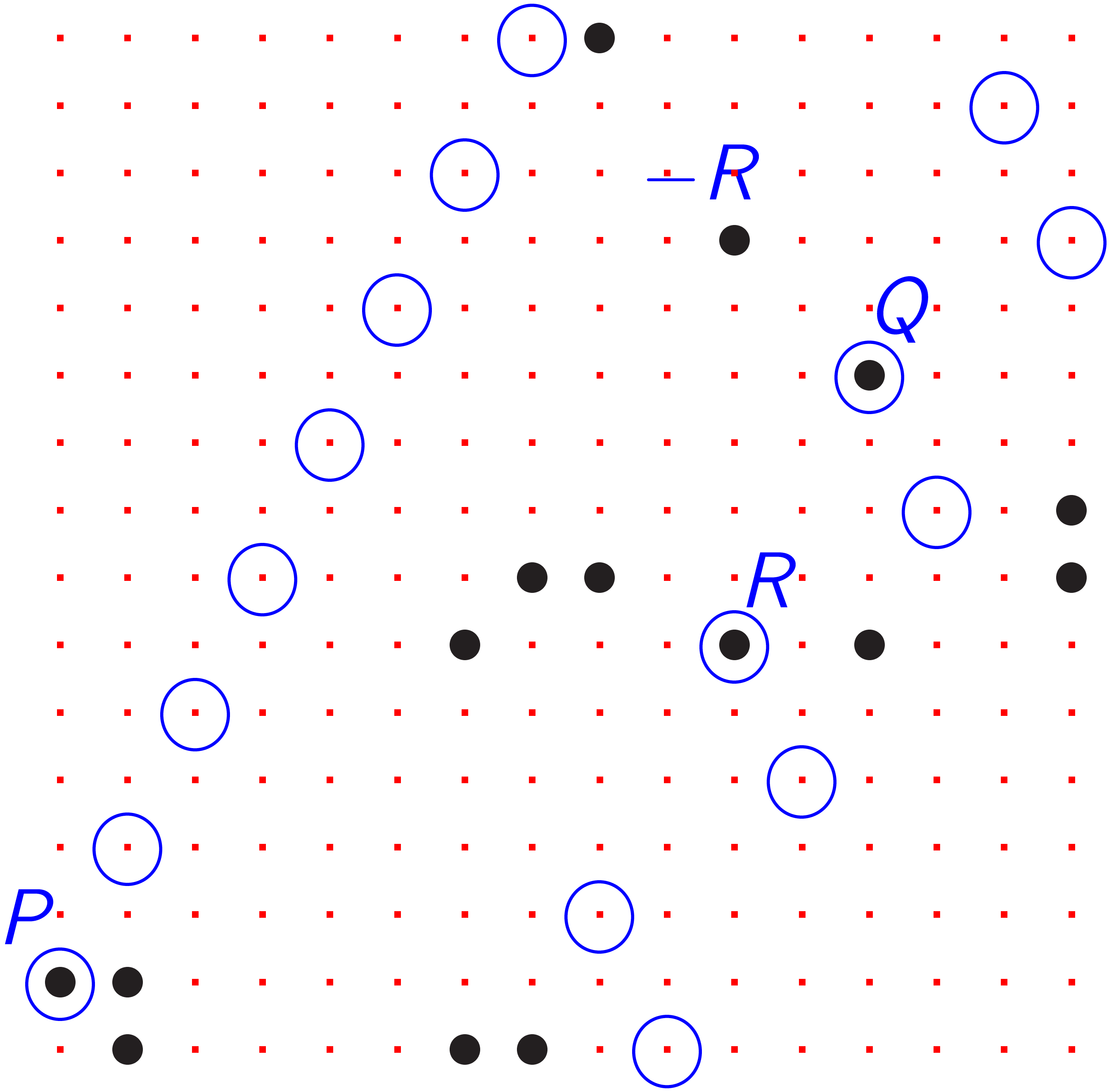
Graph of the "set of points on the elliptic curve $y^2 - 5xy = x^3 - 7$ over $(\mathbf{Z}/2)[t]/(t^4 - t - 1)$":



Line $y = tx + 1$:



$P + Q = -R$:



More elliptic curves

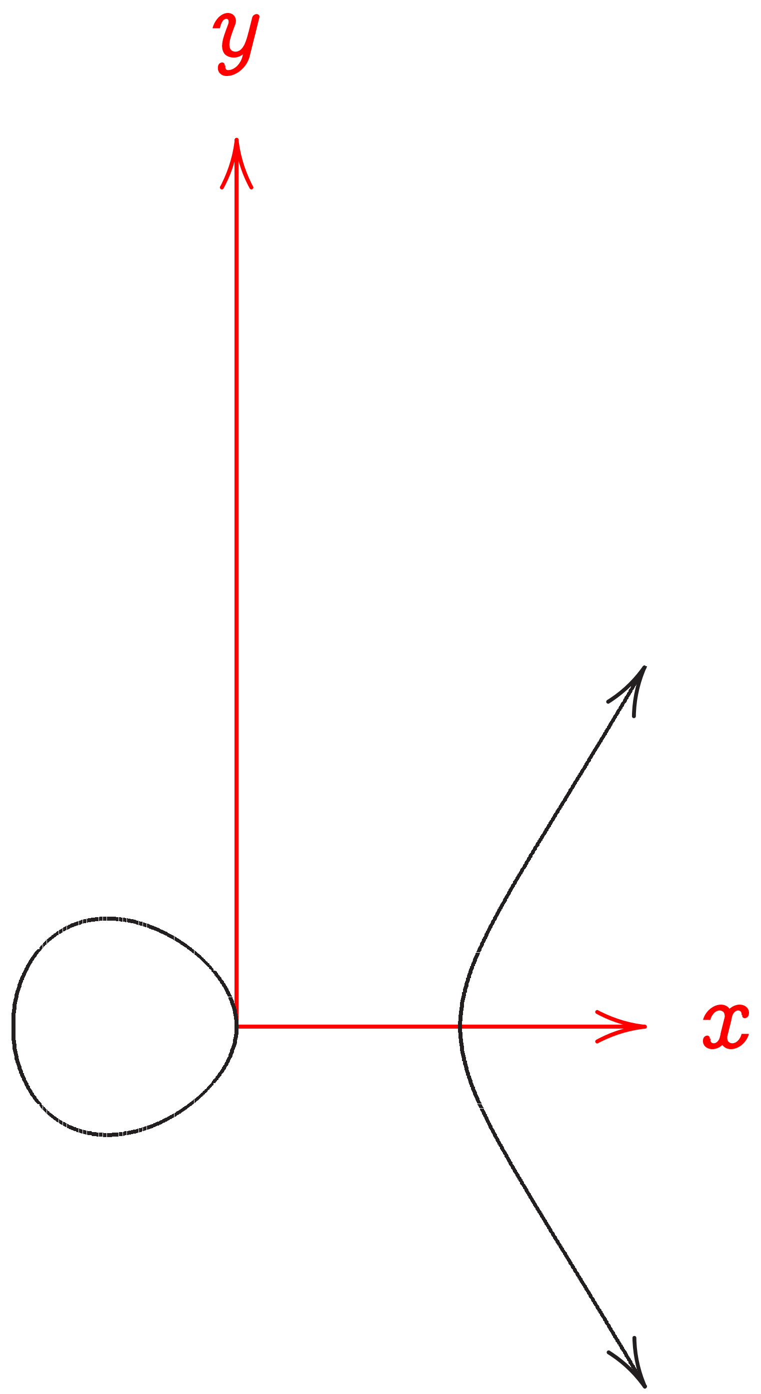
Can use any field k.

Can use any nonsingular curve $y^2 + a_1 xy + a_3 y = x^3 + a_2 x^2 + a_4 x + a_6$.

"Nonsingular": no $(x, y) \in k \times k$ simultaneously satisfies $y^2 + a_1 xy + a_3 y = x^3 + a_2 x^2 + a_4 x + a_6$ and $2y + a_1 x + a_3 = 0$ and $a_1 y = 3x^2 + 2a_2 x + a_4$.

Easy to check nonsingularity. Almost all curves are nonsingular when k is large.

e.g. $y^2 = x^3 - 30x$:



$\{(x, y) \in k \times k :$
$y^2 + a_1xy + a_3y =$
$x^3 + a_2x^2 + a_4x + a_6\} \cup \{\infty\}$
is a commutative group with standard definition of 0, $-$, $+$.
Points on line add to 0 with appropriate multiplicity.

Group is usually called "$E(k)$" where E is "the elliptic curve $y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6$."

Fairly easy to write down explicit formulas for 0, $-$, $+$ as before.

Using explicit formulas can quickly compute nth multiples in $E(k)$ given $n \in \{0, 1, 2, \ldots, 2^{256} - 1\}$ and given E, k with $\#k \approx 2^{256}$.

(How quickly?
We'll study this later.)

"Elliptic-curve discrete-logarithm problem" (ECDLP):
given points P and nP, find n.

Can find curves where ECDLP seems extremely difficult:
$\approx 2^{128}$ operations.

See “Handbook of elliptic and hyperelliptic curve cryptography” for much more information.

Two examples of elliptic curves useful for cryptography:

“NIST P-256”: $E(\mathbf{Z}/p)$ where p is the prime $2^{256}-2^{224}+2^{192}+2^{96}-1$ and E is the elliptic curve $y^2 = x^3 - 3x +$ (a particular constant).

“Curve25519”: $E(\mathbf{Z}/p)$ where p is the prime $2^{255} - 19$ and E is the elliptic curve $y^2 = x^3 + 486662x^2 + x$.

Fast arithmetic

1. Someone specifies k.
How quickly can we
perform arithmetic in k?

2. Someone specifies k and E.
How quickly can we
compute nth multiples in $E(k)$?

3. How quickly can we
compute nth multiples in $E(k)$
if we *choose* k and E?

Some examples of finite fields:

$\mathbf{Z}/(2^{255} - 19)$.
$(\mathbf{Z}/(2^{61} - 1))[t]/(t^5 - 3)$.
$(\mathbf{Z}/223))[t]/(t^{37} - 2)$.
$(\mathbf{Z}/2)[t]/(t^{283} - t^{12} - t^7 - t^5 - 1)$.

How quickly can we
add, subtract, multiply
in these fields?

Answer will depend on platform:
AMD Athlon, Sun UltraSPARC IV,
Intel 8051, Xilinx Spartan-3, etc.
Warning: different platforms
often favor different fields!

Fast integer arithmetic

How to multiply big integers?

Child's answer: Use polynomial with coefficients in $\{0, 1, \ldots, 9\}$ to represent integer in radix 10.

With this representation, multiply integers in two steps:

1. Multiply polynomials.
2. "Carry" extra digits.

Polynomial multiplication involves *small* integers.
Have split one big multiplication into many small operations.

Example of representation:

$839 = 8 \cdot 10^2 + 3 \cdot 10^1 + 9 \cdot 10^0 =$ value (at $t = 10$) of polynomial $8t^2 + 3t^1 + 9t^0$.

Squaring: $(8t^2 + 3t^1 + 9t^0)^2 =$ $64t^4 + 48t^3 + 153t^2 + 54t^1 + 81t^0$.

Carrying:

$64t^4 + 48t^3 + 153t^2 + 54t^1 + 81t^0$;

$64t^4 + 48t^3 + 153t^2 + 62t^1 + 1t^0$;

$64t^4 + 48t^3 + 159t^2 + 2t^1 + 1t^0$;

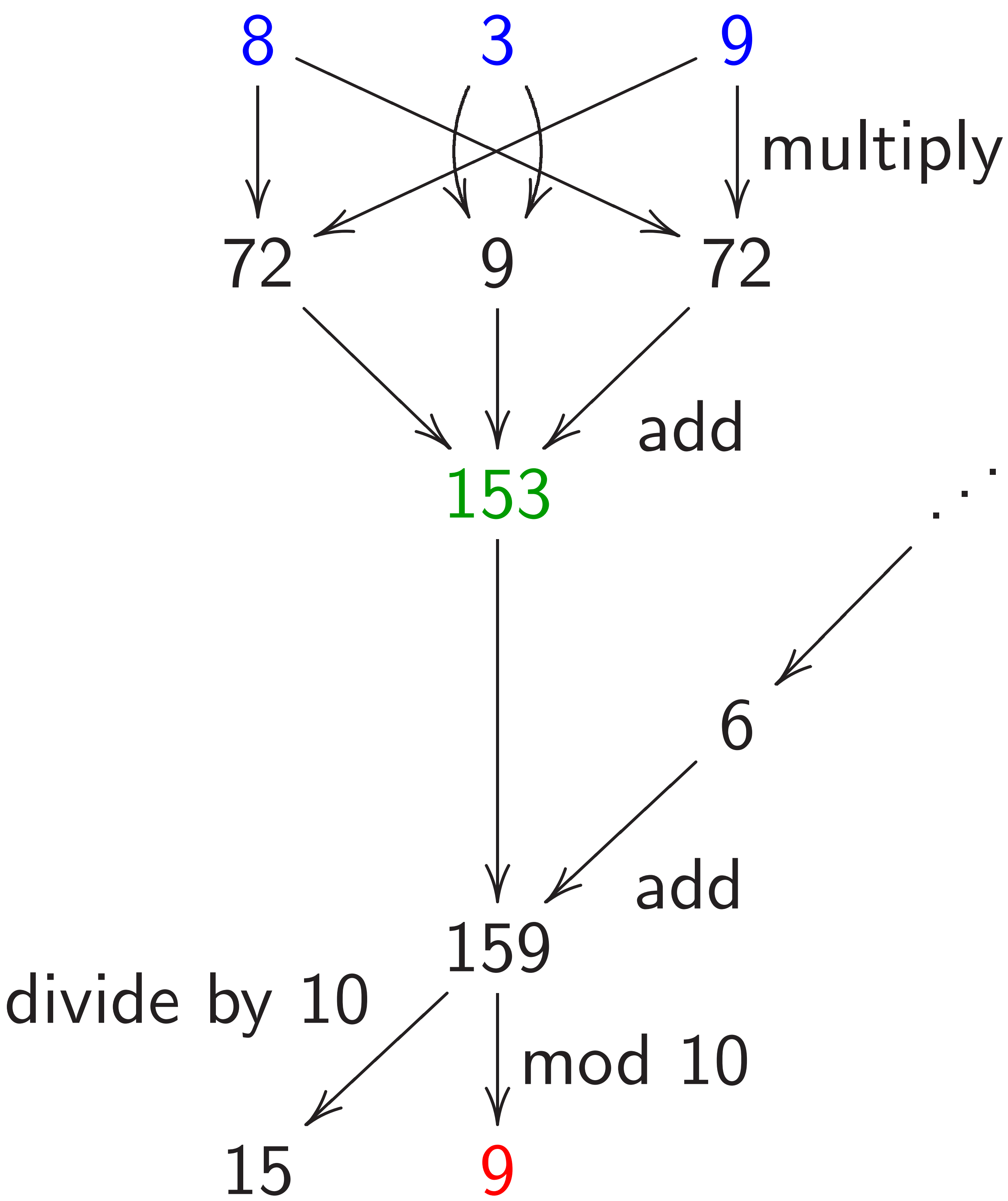
$64t^4 + 63t^3 + 9t^2 + 2t^1 + 1t^0$;

$70t^4 + 3t^3 + 9t^2 + 2t^1 + 1t^0$;

$7t^5 + 0t^4 + 3t^3 + 9t^2 + 2t^1 + 1t^0$.

In other words, $839^2 = 703921$.

What operations were used here?



Scaled variation:

$839 = 800 + 30 + 9 =$

value (at $t = 1$) of polynomial

$800t^2 + 30t^1 + 9t^0$.

Squaring: $(800t^2 + 30t^1 + 9t^0)^2 =$

$640000t^4 + 48000t^3 + 15300t^2 +$

$540t^1 + 81t^0$.

Carrying:

$640000t^4 + 48000t^3 + 15300t^2 +$

$540t^1 + 81t^0$;

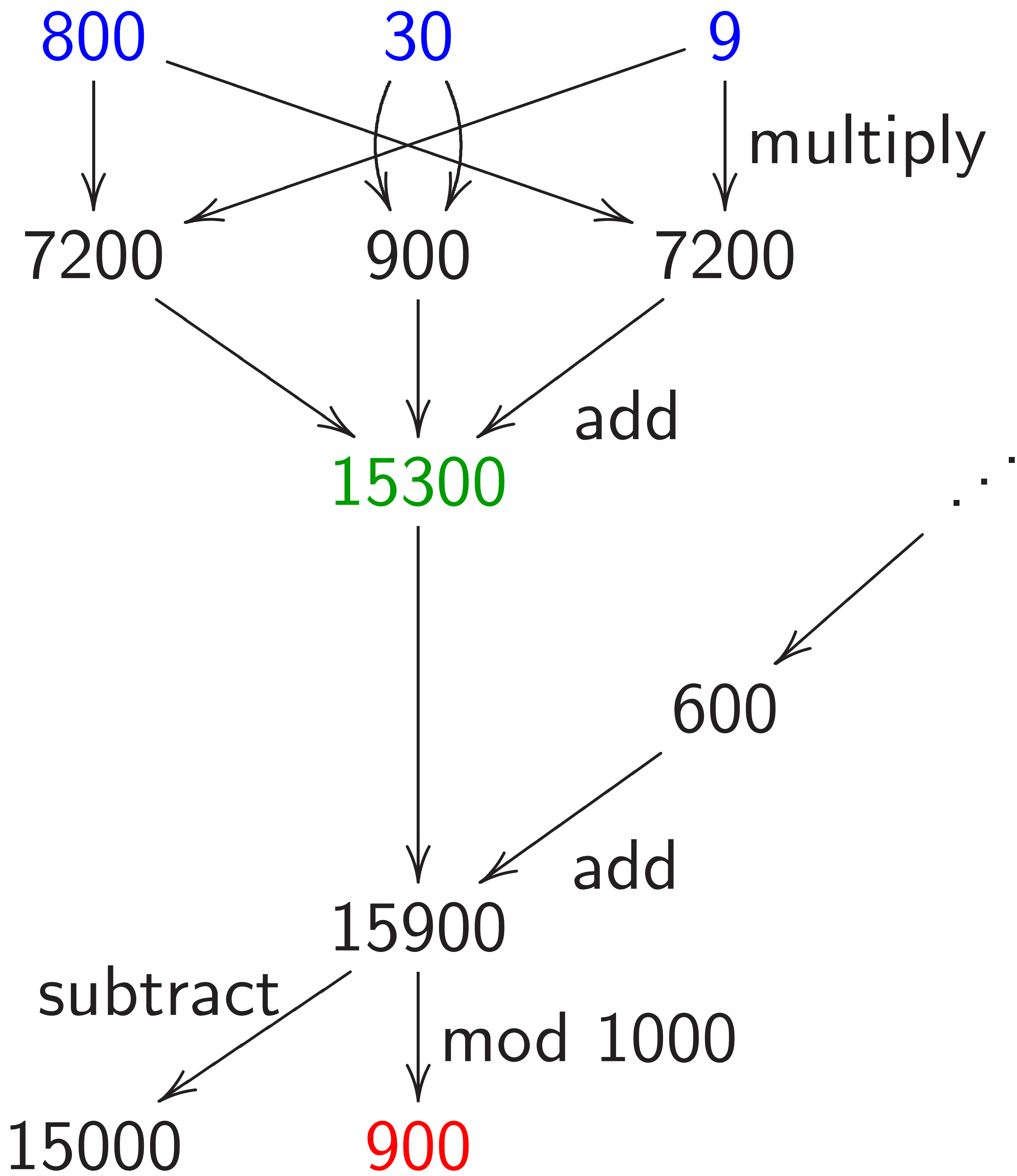
$640000t^4 + 48000t^3 + 15300t^2 +$

$620t^1 + 1t^0$; ...

$700000t^5 + 0t^4 + 3000t^3 + 900t^2 +$

$20t^1 + 1t^0$.

What operations were used here?



Speedup: double inside squaring

Squaring $\cdots + f_2 t^2 + f_1 t^1 + f_0 t^0$ produces coefficients such as $f_4 f_0 + f_3 f_1 + f_2 f_2 + f_1 f_3 + f_0 f_4$.

Compute more efficiently as $2 f_4 f_0 + 2 f_3 f_1 + f_2 f_2$.
Or, slightly faster, $2(f_4 f_0 + f_3 f_1) + f_2 f_2$.
Or, slightly faster, $(2f_4) f_0 + (2f_3) f_1 + f_2 f_2$ after precomputing $2f_1, 2f_2, \ldots$.

Have eliminated $\approx 1/2$ of the work if there are many coefficients.

Speedup: allow negative coeffs

Recall $159 \mapsto 15, 9$.

Scaled: $15900 \mapsto 15000, 900$.

Alternative: $159 \mapsto 16, -1$.

Scaled: $15900 \mapsto 16000, -100$.

Use digits $\{-5, -4, \ldots, 4, 5\}$ instead of $\{0, 1, \ldots, 9\}$.

Several small advantages: easily handle negative integers; easily handle subtraction; reduce products a bit.

Speedup: delay carries

Computing (e.g.) big $ab + c^2$: multiply a, b polynomials, carry, square c poly, carry, add, carry.

e.g. $a = 314$, $b = 271$, $c = 839$:
$(3t^2 + 1t^1 + 4t^0)(2t^2 + 7t^1 + 1t^0) =$
$6t^4 + 23t^3 + 18t^2 + 29t^1 + 4t^0$;
carry: $8t^4 + 5t^3 + 0t^2 + 9t^1 + 4t^0$.

As before $(8t^2 + 3t^1 + 9t^0)^2 =$
$64t^4 + 48t^3 + 153t^2 + 54t^1 + 81t^0$;
$7t^5 + 0t^4 + 3t^3 + 9t^2 + 2t^1 + 1t^0$.

$+$: $7t^5 + 8t^4 + 8t^3 + 9t^2 + 11t^1 + 5t^0$;
$7t^5 + 8t^4 + 9t^3 + 0t^2 + 1t^1 + 5t^0$.

Faster: multiply a, b polynomials, square c polynomial, add, carry.

$$(6t^4 + 23t^3 + 18t^2 + 29t^1 + 4t^0) + (64t^4 + 48t^3 + 153t^2 + 54t^1 + 81t^0) = 70t^4 + 71t^3 + 171t^2 + 83t^1 + 85t^0;$$

$$7t^5 + 8t^4 + 9t^3 + 0t^2 + 1t^1 + 5t^0.$$

Eliminate intermediate carries. Outweighs cost of handling slightly larger coefficients.

Important to carry between multiplications (and squarings) to reduce coefficient size; but carries are usually a bad idea for additions, subtractions, etc.

Speedup: polynomial Karatsuba

Computing product of polys f, g with (e.g.) $\deg f < 20$, $\deg g < 20$: 400 coefficient mults, 361 coefficient adds.

Faster: Write f as $F_0 + F_1 t^{10}$ with $\deg F_0 < 10$, $\deg F_1 < 10$. Similarly write g as $G_0 + G_1 t^{10}$.

Then $fg = (F_0 + F_1)(G_0 + G_1)t^{10}$ $+ (F_0 G_0 - F_1 G_1 t^{10})(1 - t^{10})$.

20 adds for $F_0 + F_1$, $G_0 + G_1$.

300 mults for three products

$F_0 G_0$, $F_1 G_1$, $(F_0 + F_1)(G_0 + G_1)$.

243 adds for those products.

9 adds for $F_0 G_0 - F_1 G_1 t^{10}$

with subs counted as adds

and with delayed negations.

19 adds for $\cdots (1 - t^{10})$.

19 adds to finish.

Total 300 mults, 310 adds.

Larger coefficients, slight expense;

still saves time.

Can apply idea recursively

as poly degree grows.

Many other algebraic speedups in polynomial multiplication: Toom, FFT, etc.

Increasingly important as polynomial degree grows. $O(n \lg n \lg\lg n)$ coeff operations to compute n-coeff product.

Useful for sizes of n that occur in cryptography? Maybe; active research area.

Using CPU's integer instructions

Replace radix 10 with, e.g., 2^{24}. Power of 2 simplifies carries.

Adapt radix to platform.

e.g. Every 2 cycles, Athlon 64 can compute a 128-bit product of two 64-bit integers. (5-cycle latency; parallelize!) Also low cost for 128-bit add.

Reasonable to use radix 2^{60}. Sum of many products of digits fits comfortably below 2^{128}. Be careful: analyze largest sum.

e.g. In 4 cycles, Intel 8051 can compute a 16-bit product of two 8-bit integers.
Could use radix 2^6.
Could use radix 2^8, with 24-bit sums.

e.g. Every 2 cycles, Pentium 4 F3 can compute a 64-bit product of two 32-bit integers.
(11-cycle latency; yikes!)
Reasonable to use radix 2^{28}.

Warning: Multiply instructions are very slow on some CPUs.
e.g. Pentium 4 F2: 10 cycles!

Using floating-point instructions

Big CPUs have separate floating-point instructions, aimed at numerical simulation but useful for cryptography.

In my experience, floating-point instructions support faster multiplication (often much, much faster) than integer instructions, except on the Athlon 64. Other advantages: portability; easily scaled coefficients.

e.g. Every 2 cycles, Pentium III can compute a 64-bit product of two floating-point numbers, and an independent 64-bit sum.

e.g. Every cycle, Athlon can compute a 64-bit product and an independent 64-bit sum.

e.g. Every cycle, UltraSPARC III can compute a 53-bit product and an independent 53-bit sum. Reasonable to use radix 2^{24}.

e.g. Pentium 4 can do the same using SSE2 instructions.

How to do carries in floating-point registers? (No CPU carry instruction: not useful for simulations.)

Exploit floating-point rounding: add big constant, subtract same constant.

e.g. Given α with $|\alpha| \le 2^{75}$: compute 53-bit floating-point sum of α and constant $3 \cdot 2^{75}$, obtaining a multiple of 2^{24}; subtract $3 \cdot 2^{75}$ from result, obtaining multiple of 2^{24} nearest α; subtract from α.

Reducing modulo a prime

Fix a prime p.
The prime field $\mathbf{Z}/p$
is the set $\{0, 1, 2, \ldots, p-1\}$
with $-$ defined as $-$ mod p,
$+$ defined as $+$ mod p,
$\cdot$ defined as $\cdot$ mod p.

e.g. $p = 1000003$:
$1000000 + 50 = 47$ in $\mathbf{Z}/p$;
$-1 = 1000002$ in $\mathbf{Z}/p$;
$117505 \cdot 23131 = 1$ in $\mathbf{Z}/p$.

How to multiply in $\mathbf{Z}/p$?

Can use definition:
$fg \bmod p = fg - p \lfloor fg/p \rfloor$.
Can multiply fg by a precomputed $1/p$ approximation; easily adjust to obtain $\lfloor fg/p \rfloor$.
Slight speedup: "2-adic inverse"; "Montgomery reduction."

We can do better: normally p is chosen with a special form (or dividing a special form; see "redundant representations") to make $fg \bmod p$ much faster.

e.g. In $\mathbf{Z}/1000003$:

$314159265358 =$

$314159 \cdot 1000000 + 265358 =$

$314159(-3) + 265358 =$

$-942477 + 265358 =$

$-677119.$

Easily adjust to range $\{0, 1, \ldots, p-1\}$ by adding/subtracting a few p's. (Beware timing attacks!)

Speedup: Delay the adjustment; extra p's won't damage subsequent field operations.

Can delay carries until after multiplication by 3.

e.g. To square 314159 in $\mathbf{Z}/1000003$: Square poly $3t^5 + 1t^4 + 4t^3 + 1t^2 + 5t^1 + 9t^0$, obtaining $9t^{10} + 6t^9 + 25t^8 + 14t^7 + 48t^6 + 72t^5 + 59t^4 + 82t^3 + 43t^2 + 90t^1 + 81t^0$.

Reduce: replace $(c_i)t^{6+i}$ by $(-3c_i)t^i$, obtaining $72t^5 + 32t^4 + 64t^3 - 32t^2 + 48t^1 - 63t^0$.

Carry: $8t^6 - 4t^5 - 2t^4 + 1t^3 + 2t^2 + 2t^1 - 3t^0$.

To minimize poly degree, mix reduction and carrying, carrying the top sooner.

e.g. Start from square $9t^{10} + 6t^9 + 25t^8 + 14t^7 + 48t^6 + 72t^5 + 59t^4 + 82t^3 + 43t^2 + 90t^1 + 81t^0$.

Reduce $t^{10} \to t^4$ and carry $t^4 \to t^5 \to t^6$: $6t^9 + 25t^8 + 14t^7 + 56t^6 - 5t^5 + 2t^4 + 82t^3 + 43t^2 + 90t^1 + 81t^0$.

Finish reduction: $-5t^5 + 2t^4 + 64t^3 - 32t^2 + 48t^1 - 87t^0$. Carry $t^0 \to t^1 \to t^2 \to t^3 \to t^4 \to t^5$: $-4t^5 - 2t^4 + 1t^3 + 2t^2 - 1t^1 + 3t^0$.

Speedup: non-integer radix

Consider $\mathbf{Z}/(2^{61}-1)$.

Five coeffs in radix 2^{13}?

$f_4 t^4 + f_3 t^3 + f_2 t^2 + f_1 t^1 + f_0 t^0$.

Most coeffs could be 2^{12}.

Square $\cdots + 2(f_4 f_1 + f_3 f_2) t^5 + \cdots$.

Coeff of t^5 could be $> 2^{25}$.

Reduce: $2^{65} = 2^4$ in $\mathbf{Z}/(2^{61}-1)$;

$\cdots + (2^5(f_4 f_1 + f_3 f_2) + f_0^2) t^0$.

Coeff could be $> 2^{29}$.

Very little room for additions, delayed carries, etc. on 32-bit platforms.

Scaled: Evaluate at $t = 1$.

f_4 is multiple of 2^{52};

f_3 is multiple of 2^{39};

f_2 is multiple of 2^{26};

f_1 is multiple of 2^{13};

f_0 is multiple of 2^{0}. Reduce:

$\cdots + (2^{-60}(f_4 f_1 + f_3 f_2) + f_0^2)t^0$.

Better: Non-integer radix $2^{12.2}$.

f_4 is multiple of 2^{49};

f_3 is multiple of 2^{37};

f_2 is multiple of 2^{25};

f_1 is multiple of 2^{13};

f_0 is multiple of 2^{0}.

Saves a few bits in coeffs.

More finite fields

Fix a prime p. Fix a poly φ in one variable t with φ irreducible mod p.

The finite field $(\mathbf{Z}/p)[t]/\varphi$ is the set of polynomials $f_{\deg\varphi-1}t^{\deg\varphi-1} + \cdots + f_1 t^1 + f_0 t^0$ with each $f_i \in \mathbf{Z}/p$ and with $-, +, \cdot$ defined modulo p and modulo φ.

$(\mathbf{Z}/p)[t]/\varphi$ is an "extension" of the prime field $\mathbf{Z}/p$; it has "characteristic" p.

e.g. 223 is prime, and poly $t^6 - 3$ is irreducible mod 223, so $(\mathbf{Z}/223)[t]/(t^6 - 3)$ is a field.

223^6 elements of field, namely polynomials $f_5 t^5 + f_4 t^4 + f_3 t^3 + f_2 t^2 + f_1 t^1 + f_0 t^0$ with each $f_i \in \{0, 1, \ldots, 222\}$.

After adding, subtracting, multiplying: replace t^6 by 3, replace t^7 by $3t$, etc.; and reduce coefficients modulo 223. e.g. $(9t^4 + 1)^2 = 81t^8 + 18t^4 + 1 = 243t^2 + 18t^4 + 1 = 18t^4 + 20t^2 + 1$.

Have two levels of polynomials when p is large: element of $(\mathbf{Z}/p)[t]/\varphi$ is poly mod φ; each poly coefficient is integer represented as poly in some radix.

e.g. $f_4t^4 + f_3t^3 + f_2t^2 + f_1t^1 + f_0t^0$ in $(\mathbf{Z}/(2^{61} - 1))[t]/(t^5 - 3)$ could have each coefficient f_i represented as poly of degree < 3 in radix $2^{61/3}$.

When p is small, especially $p = 2$, many speedups beyond this talk: batching coefficients, using fast Frobenius, et al.