GUARANTEED MESSAGE AUTHENTICATION FASTER THAN MD5 (ABSTRACT)

DANIEL J. BERNSTEIN

Let $r_0, r_1, r_2, r_3, x_0, x_1, x_2, x_3, m_0, m_1, m_2, \ldots, m_{n-1}$ be integers in $[-2^{31}, 2^{31}-1]$. Define $r = 2^{96}r_3 + 2^{64}r_2 + 2^{32}r_1 + r_0$, $x = 2^{96}x_3 + 2^{64}x_2 + 2^{32}x_1 + x_0$, and

$$s = (r^{n+1} + r^n m_0 + r^{n-1} m_1 + \cdots + r m_{n-1} + x) \bmod (2^{127} - 1).$$

I can compute s in about $330 + 19n$ Pentium cycles, or $470 + 26n$ UltraSPARC cycles, after a precomputation depending only on r.

Applications. Here's one way to mathematically guarantee the authenticity of a single message $m = (m_0, m_1, \ldots, m_{n-1})$. The sender and receiver share a secret uniform random pair (r, x). The sender computes s as above and sends (m, s). The receiver verifies (m, s) by recomputing s. An attacker, given (m, s), has a negligible probability of successfully forging a different message.

This system is faster than yesterday's MD5-based systems: $s = \mathrm{MD5}(r, m, x)$, for example, or $s = \mathrm{MD5}(r, \mathrm{MD5}(x, m))$. It provides essentially the same level of security against today's attacks; unlike MD5, it is also guaranteed secure against tomorrow's attacks. It can easily be extended to handle multiple messages.

The same method can be used to reduce a multiprecision integer modulo a big secret prime. One application is to check ring equations involving large integers—for example, the equation $s^2 = tn + fh$ in my variant of the Rabin-Williams public-key signature system—with a negligible chance of error.

Method. The following comments apply to the Pentium.

The fastest way to add and multiply small integers is with the help of the floating-point unit. For example, it takes just one cycle to compute an exact product of two 32-bit integers in floating-point registers. There is an old trick to split the result into low bits and high bits with a few floating-point additions and subtractions; one need not retrieve integers from the floating-point unit.

In multiprecision arithmetic one should generally use a radix below 2^{32} so that more useful work can be done between carries. I precompute small integers $c_{i,j}$ such that r^i is congruent to $c_{i,0} + 2^{26}c_{i,1} + 2^{52}c_{i,2} + 2^{78}c_{i,3} + 2^{104}c_{i,4}$ modulo $2^{127} - 1$. If n is not too large then the dot products $\sum_i c_{n-i,0}m_i$, $\sum_i c_{n-i,1}m_i$, etc. fit safely into floating-point registers. Handling large n is not much more difficult.

The `gcc -O6` optimizer does a poor job of instruction scheduling and register allocation. I use `gcc -O1` and schedule instructions manually. The speeds reported above are still not optimal.

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History. There are many previous message authentication systems that reduce various rings modulo big secret random prime ideals. This is the first high-security system to break the MD5 speed barrier. The direct ancestor of this system is Shoup's analogous system in characteristic 2.

Department of Mathematics, Statistics, and Computer Science (M/C 249), The University of Illinois at Chicago, Chicago, IL 60607–7045

E-mail address: djb@pobox.com